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## THE DOUBLE POWERLOCALE AND EXPONENTIATION: A CASE STUDY IN GEOMETRIC LOGIC

STEVEN VICKERS

ABSTRACT. If  $X$  is a locale, then its double powerlocale  $\mathbb{P}X$  is defined to be  $P_U(P_L(X))$  where  $P_U$  and  $P_L$  are the upper and lower powerlocale constructions. We prove various results relating it to exponentiation of locales, including the following. First, if  $X$  is a locale for which the exponential  $\mathbb{S}^X$  exists (where  $\mathbb{S}$  is the Sierpinski locale), then  $\mathbb{P}X$  is an exponential  $\mathbb{S}^{\mathbb{S}^X}$ . Second, if in addition  $W$  is a locale for which  $\mathbb{P}W$  is homeomorphic to  $\mathbb{S}^X$ , then  $X$  is an exponential  $\mathbb{S}^W$ .

The work uses geometric reasoning, i.e. reasoning stable under pullback along geometric morphisms, and this enables the locales to be discussed in terms of their points as though they were spaces. It relies on a number of geometricity results including those for locale presentations and for powerlocales.

### 1. Introduction

Why should exponentiable locales be locally compact? That is, if a locale  $X$  is such that the exponential  $Y^X$  exists for every locale  $Y$ , why does it follow that the frame  $\Omega X$  of opens is a continuous lattice? The proof in [Johnstone 82] involves an analysis of injective locales, but [Hyland 81] gives a more direct proof. We paraphrase it in a naive way that, on the face of it, is just too simple to be right.

1. If  $X$  is exponentiable, then, in particular,  $\mathbb{S}^X$  exists, where  $\mathbb{S}$  is the *Sierpinski* locale whose frame is free on one generator. The points of  $\mathbb{S}^X$  are in bijective (and order preserving – cf. Hyland’s sublemma 2.1) correspondence with the continuous maps from  $X$  to  $\mathbb{S}$ , and hence with the opens of  $X$ . We should therefore like to think of  $\mathbb{S}^X$  as a particular topology on the frame  $\Omega X$ , and one whose specialization order is the frame order.
2. The evaluation map  $\text{ev} : \mathbb{S}^X \times X \rightarrow \mathbb{S}$  corresponds to an open of  $\mathbb{S}^X \times X$ , and that can be expressed in the form

$$\text{ev} = \bigvee \{V \times b \mid V \in \Omega \mathbb{S}^X, b \in \Omega X, V \times b \leq \text{ev}\}.$$

3. If  $V \times b \leq \text{ev}$  then  $b$  is a lower bound of  $V$ . This is because if  $a$  is in  $V$  (we are thinking of  $V$  as an open subset of  $\Omega X$  here) and  $x$  is in  $b$  then  $\langle a, x \rangle$  is in  $\text{ev}$  and

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so  $x$  is in  $a$ . Hence  $b \leq a$ . Indeed,  $b \ll a$  ( $b$  is way below  $a$ ), for suppose  $a \leq \bigvee_i^\uparrow c_i$ . In a locale every open  $V$  is also Scott open, and it follows that  $c_i \in V$  for some  $i$ , and then  $b \leq c_i$ . (cf. Hyland’s sublemma 2.3.)

4. If  $a$  is an open of  $X$ , then (Hyland’s sublemma 2.4)

$$a = \bigvee^\uparrow \{b \in \Omega X \mid \exists V \in \Omega \mathbb{S}^X. (V \times b \leq \text{ev}, a \in V)\}$$

This join is directed, for if  $V_i \times b_i \leq \text{ev}$  and  $a \in V_i$  ( $i = 1, 2$ ), then  $a \in V_1 \wedge V_2$ ,  $(V_1 \wedge V_2) \times b_i \leq \text{ev}$  and so  $(V_1 \wedge V_2) \times (b_1 \vee b_2) \leq \text{ev}$ . Also, the set of joined elements is inhabited, for it contains  $b = \emptyset$  because  $\mathbb{S}^X \times \emptyset = \emptyset \leq \text{ev}$ .

To show that  $a$  is this join, suppose  $x \in a$ . Then  $\langle a, x \rangle \in \text{ev}$  and so  $\langle a, x \rangle \in V \times b \leq \text{ev}$  for some  $V, b$ .

5. Hence every open of  $X$  is a join of opens way below it, which suffices to show that  $\Omega X$  is a continuous lattice.

We have used very spatial modes of reasoning here, treating opens as sets of points. On the other hand, we have also used the *sobriety* of locales in step 3, since that is what tells us that every open is Scott open. It appears to be an argument that could only work for sober spaces or spatial locales, and if we tried to restrict it to those we should have the problem of saying just which category  $X$  is exponentiable in.

The flaw in such spatial arguments is, of course, that in general locales are not spatial: they do not have enough points. It is for this reason that one usually approaches locales through pure lattice theory, as captured in the phrase “point-free topology” [Johnstone 82].

An underlying goal of this paper is to show that, nonetheless, naive spatial arguments of this kind can in fact be essentially correct.

The correctness does not come for free, however. If the points of  $X$  are understood as the maps from  $1$  to  $X$ , in other words the *global* points, then there really are not enough in general. It is only by understanding points in a *generalized* fashion, maps from an arbitrary locale  $Y$  to  $X$ , that one finds enough. Reasoning about these points can be carried through by replacing ordinary sets by sheaves over  $Y$ , but these do not conform with all classical reasoning principles – one must instead reason constructively.

This, then, is a practical benefit of constructive reasoning. If it is of the kind that is valid in toposes (which we shall refer to as *intuitionistic* reasoning), it can be applied to give results about sheaves, treating them as though they were sets. This is perhaps most strikingly displayed in [JoyTie 84]. This also allows results about locales to be turned for free into results about continuous maps between locales, since a map targeted on  $X$  is equivalent to an internal locale in the category  $\mathcal{S}X$  of sheaves over  $X$ .

A further refinement is to use the more stringent *geometric* reasoning, stable under pullback along geometric morphisms. Under this regime, reasoning in  $\mathcal{S}X$  about the

generic point of  $X$  can be transferred to *all* points of  $X$ , and this can be exploited to give a very spatial view of locale theory. A locale is then defined by saying (geometrically) what its points are, and a geometric morphism – a *map* – is defined by saying (geometrically) how it transforms points. Topology is intrinsic, and continuity is automatic. This technique, actually an old trade secret of topos theorists, is used very explicitly in [Vickers 99], [Vickers 01].

The present paper develops this in an example where geometric constructions must be used not only on points, but even on locales themselves. To see this in overview, the main result is that if a locale  $X$  is exponentiable then so is  $\mathbb{S}^X$  (where  $\mathbb{S}$  is the Sierpinski locale), and  $\mathbb{S}^{\mathbb{S}^X}$  is homeomorphic to a different construction  $\mathbb{P}X$ , the *double powerlocale* of the title. (It is defined as  $P_U P_L X \cong P_L P_U X$ , but we shall not use that for the moment.)

Looking at global points, this implies an equivalence between points of  $\mathbb{P}X$  and opens of  $\mathbb{S}^X$  (i.e. maps  $\mathbb{S}^X \rightarrow \mathbb{S}$ ). However, the definition of exponential goes beyond this to require an equivalence between maps  $Y \times \mathbb{S}^X \rightarrow \mathbb{S}$  and maps  $Y \rightarrow \mathbb{P}X$  for any locale  $Y$ . Consider what happens when we work intuitionistically in  $\mathcal{S}Y$ . The maps  $Y \times \mathbb{S}^X \rightarrow \mathbb{S}$  are equivalent to opens of a locale over  $Y$  given by the projection  $\text{fst} : Y \times \mathbb{S}^X \rightarrow Y$ . If we have a proof for global points that is intuitionistic then we can apply it in  $\mathcal{S}Y$  to get an equivalence with points over  $Y$  of  $\mathbb{P}_Y(Y \times X \rightarrow Y)$  where  $\mathbb{P}_Y$  is the  $\mathbb{P}$  construction carried out in  $\mathcal{S}Y$ .

$$\begin{array}{ccc}
 Y \times X & & \mathbb{P}_Y(Y \times X \rightarrow Y) \\
 \searrow & & \swarrow \\
 & Y & \\
 \text{fst} & & p
 \end{array}$$

Such a point over  $Y$  is a map  $e : Y \rightarrow \mathbb{P}_Y(Y \times X \rightarrow Y)$  such that  $e; p = \text{Id}_Y$ . However, it remains to make a connection between  $\mathbb{P}_Y$  and  $\mathbb{P}$ , and what is needed is that  $\mathbb{P}$  itself should be geometric, preserved by pullback. This is to say that

$$\mathbb{P}_Y(Y \times X \rightarrow Y) \cong (Y \times \mathbb{P}X \rightarrow Y)$$

and then points of  $Y \times \mathbb{P}X$  over  $Y$  are just maps  $Y \rightarrow \mathbb{P}X$ .

The geometricity of a construction such as  $\mathbb{P}$  is highly non-trivial. This is because locales are defined as frames (certain complete lattices), but the frames are *not* geometric. If  $f : X \rightarrow Y$  is a map and  $A$  is a frame in  $\mathcal{S}Y$ , then  $f^*(A)$  is not in general a frame in  $\mathcal{S}X$ . The reason is that completeness of  $A$  in  $\mathcal{S}Y$  means having all  $Y$ -set-indexed joins, joins of all families  $(a_i)_{i \in I}$  indexed by objects  $I$  of  $\mathcal{S}Y$ . This property is preserved in  $f^*(A)$ , but to be a frame in  $\mathcal{S}X$  it needs more general  $X$ -set-indexed joins. In effect,  $f^*(A)$  is only a basis, not the full topology, for the pulled back locale, and this provides the essential idea for our geometric working. We must deal not directly with the frames themselves but with presentations of them (by generators and relations) and show how  $\mathbb{P}$  can be constructed geometrically in terms of the presentations. In fact this result is one of considerable depth, relying on “coverage theorems” to relate presentations of frames to those of suplattices and preframes, and showing that some messy geometric constructions can be given universal characterizations intuitionistically.

That, then, is the underlying shape of the paper: geometricity results can be exploited to simplify some localic reasoning. We now move on to discussing the double powerlocale  $\mathbb{P}$ .

In locale theory one encounters a number of ways of forming locales whose points are sublocales of other locales. Obvious examples are the powerlocales (upper, lower and Vietoris), and in addition if a locale  $X$  is exponentiable then the points of  $\mathbb{S}^X$  (where  $\mathbb{S}$  is the Sierpinski locale) are the maps from  $X$  to  $\mathbb{S}$  and hence in bijective correspondence with the open or with the closed sublocales of  $X$ .

Some interesting ways in which these relate to each other are shown in the case of algebraic dcpos.

1.1. PROPOSITION. *If  $X$  is an algebraic dcpo then  $\mathbb{S}^{\mathbb{S}^X} \cong P_L(P_U(X))$  where  $P_L$  and  $P_U$  denote the lower and upper powerlocales.*

PROOF. An algebraic dcpo is one of the form  $\text{Idl } P$  where  $P$  is a poset. Its frame of Scott opens is isomorphic to the Alexandrov topology on  $P$ , an open being an upper closed subset of  $P$ . This is equivalent to a filter of  $\mathcal{F}P$  (the finite powerset) under the upper preorder

$$S \sqsubseteq_U T \text{ iff } \forall t \in T. \exists s \in S. s \sqsubseteq t$$

This is the same as an ideal of  $\mathcal{F}P$  under  $(\sqsubseteq_U)^{\text{op}}$ , and in fact  $\text{Idl}(P)$  is exponentiable with  $\mathbb{S}^X$  homeomorphic to the ideal completion of  $(\mathcal{F}P, (\sqsubseteq_U)^{\text{op}})$ . On the other hand, the upper powerlocale is the ideal completion of  $(\mathcal{F}P, \sqsubseteq_U)$ , so for algebraic dcpos  $X$  the exponential  $\mathbb{S}^X$  is found by first taking the upper powerlocale and then taking the ‘‘Hofmann-Lawson’’ dual that inverts the generating poset (of compact points). A consequence is that  $\mathbb{S}^{\mathbb{S}^X}$  is got by first taking the upper powerlocale, then dualizing, then taking the upper powerlocale again, then dualizing again. But the last three steps are equivalent to taking the lower powerlocale, for that is given by the ideal completion of  $(\mathcal{F}P, \sqsubseteq_L)$  where

$$S \sqsubseteq_L T \text{ iff } \forall s \in S. \exists t \in T. s \sqsubseteq t$$

Hence for an algebraic dcpo  $X$  we have  $\mathbb{S}^{\mathbb{S}^X} \cong P_L(P_U(X))$ , which we write as  $\mathbb{P}X$ . ■

We shall show (Theorem 11.1) that this result extends to *all* exponentiable locales.

To some extent this is no surprise. It is a result stating that the points of  $\mathbb{P}X$  are in bijective correspondence with the maps from  $\mathbb{S}^X$  to  $\mathbb{S}$ , in other words the opens of  $\mathbb{S}^X$ . By [JoVic 91], the global points of  $\mathbb{P}X$  are known to be the functions from  $\Omega X$  to  $\Omega$  that preserve directed joins. On the other hand, it is also known that if  $\mathbb{S}^X$  exists then  $X$  is exponentiable and  $\Omega X$  is a continuous lattice whose Scott topology gives the locale  $\mathbb{S}^X$ . Hence the opens of  $\mathbb{S}^X$ , the Scott opens of  $\Omega X$ , are just the Scott continuous maps from  $\Omega X$  to  $\Omega$  – in other words, the functions that preserve directed joins.

We therefore know already that the global points of  $\mathbb{P}X$  correspond to the global opens of  $\mathbb{S}^X$ . However, we do not rely on established results about exponentiability, but show direct from the definition that if  $\mathbb{S}^X$  exists then  $\mathbb{S}^{\mathbb{S}^X}$  (exists too and) is homeomorphic to  $\mathbb{P}X$ . In a partial converse, we also show that if  $\mathbb{S}^X$  is homeomorphic to  $\mathbb{P}W$  then  $\mathbb{S}^W$

exists and is homeomorphic to  $X$ . In the process we also prove many properties of  $\mathbb{P}$  and  $\mathbb{P}$ -algebras.

Recent work of Paul Taylor [Taylor 01] has investigated an abstract formalism that includes the monad  $\Sigma^2 : X \mapsto \mathbb{S}^{\mathbb{S}^X}$  on locally compact (i.e. exponentiable) locales (the Sierpinski locale is  $\Sigma$  in Taylor's notation). The  $\Sigma^2$ -algebras are frames, and amongst other topics Taylor considers the  $\Sigma^2$ -algebra structure on  $\mathbb{S}$  and (hence) on exponentials  $\mathbb{S}^X$ , showing for instance a Stone duality result that maps  $X \rightarrow Y$  are equivalent to  $\Sigma^2$ -homomorphisms  $\mathbb{S}^Y \rightarrow \mathbb{S}^X$ .

Another purpose therefore of the present paper is to show how  $\Sigma^2$  is the restriction to locally compact locales of a monad ( $\mathbb{P}$ ) that does not rely on exponentiability.

## 2. Miscellaneous remarks on locales and toposes

**LOCALES AND MAPS.** We use the words “locale” and “frame” following [Johnstone 82] (as opposed to “space” and “locale” in [JoyTie 84]) and we shall always take care to use notation and language for locales that is appropriate to their spatial aspect. We write **Loc** for the category of locales. If  $X$  is a locale, then we write  $\Omega X$  for its corresponding frame of opens and  $\mathcal{S}X$  for its category of sheaves. We write **true** and **false** for the top and bottom elements of a frame.

If  $X$  and  $Y$  are locales, then notation such as  $f : X \rightarrow Y$  will always denote a continuous map from  $X$  to  $Y$ , corresponding to a frame homomorphism  $\Omega f$  (also written as  $f^*$ ) from  $\Omega Y$  to  $\Omega X$ .

**DCPO ENRICHMENT.** We write  $\sqsubseteq$  for the specialization order on **Loc**,

$$f \sqsubseteq g \text{ iff } \forall b \in \Omega Y. \Omega f(b) \leq \Omega g(b)$$

**Loc** is dcpo (directed complete poset) enriched under  $\sqsubseteq$ , with directed joins defined by

$$\Omega(\bigsqcup_i^\uparrow f_i)(b) = \bigvee_i^\uparrow \Omega f_i(b)$$

(This will be elaborated in Section 6.) The superscript  $\uparrow$  after a join symbol will indicate that the join is directed.

**PRESENTATIONS.** Crucial to our development will be the use of *generators and relations* for *presenting* frames (see in particular [Vickers 89]), using notation such as

$$\text{Fr}\langle G|R\rangle$$

Here  $G$  is a set of generators and  $R$  a set of relations of the form  $\phi_1 \leq \phi_2$  or  $\phi_1 = \phi_2$ ,  $\phi_1$  and  $\phi_2$  being expressions formed from the generators using finite meets and arbitrary joins. (This will be formalized in Section 5.) If  $G$  has structure (e.g. poset or semilattice) that is to be preserved in the frame, then it is convenient to indicate this by the word “qua”, for instance, “ $G$  (qua  $\wedge$ -semilattice)”.

SIERPINSKI LOCALE. The *Sierpinski locale*  $\mathbb{S}$  is defined by

$$\Omega\mathbb{S} = \text{Fr}\langle 1 \rangle,$$

i.e. the free frame on one generator. Its points are the subsets of  $1$ , and in particular we have two classical points  $\perp = \emptyset$  and  $\top = 1$ . We write  $\{\top\}$  for the generating open, since  $\top$  is the sole point of  $\mathbb{S}$  in it.

From the definition it is clear that there is a bijection between opens of  $X$  and maps  $X \rightarrow \mathbb{S}$ . By abuse of notation, we shall use the same symbol  $a$  for –

- an open  $a \in \Omega X$ ,
- the map  $a : X \rightarrow \mathbb{S}$ ,
- the open sublocale  $a \hookrightarrow X$ .

However, if  $X$  is exponentiable, we shall use different notation  $\ulcorner a \urcorner$  for the corresponding point of  $\mathbb{S}^X$ .

TOPOSES. From time to time, mostly in discussing the constructive reasoning, we shall refer to toposes. Just as with locales, we shall use notation and language adapted to their aspect as generalized spaces rather than generalized universes of sets. For a topos  $X$ , we write  $\mathcal{S}X$  for its category of sheaves, its generalized universe of sets. This is its ordinary category-theoretic expression: in the conventional language one simply says that the topos *is* this category. Note also that what we refer to as sheaves over a topos are exactly the objects of this category. We also write  $\Omega X$  for its frame of opens, the frame of subobjects of  $1$  (or global elements of  $\Omega$ ) in  $\mathcal{S}X$ . Notationally a locale  $X$  is thus just a special “localic” kind of topos, one in which the whole of  $\mathcal{S}X$  is determined by  $\Omega X$ . We make no attempt to say whether a locale “is” its frame or “is” its category of sheaves. It simply *has* both and can be defined by either.

If  $X$  and  $Y$  are toposes, then notation such as  $f : X \rightarrow Y$  will always denote a geometric morphism from  $X$  to  $Y$ , never a functor from  $\mathcal{S}X$  to  $\mathcal{S}Y$ . We shall often refer to this as a *map* from  $X$  to  $Y$ .

Note that the inverse image functor  $f^* : \mathcal{S}Y \rightarrow \mathcal{S}X$  is pullback along  $f$  in the following sense. Sheaves over  $Y$  (objects of  $\mathcal{S}Y$ ) correspond, up to homeomorphism, with local homeomorphisms (étale maps) targeted at  $Y$ , and then  $f^*$  on sheaves corresponds to pullback along  $f$  (also written  $f^*$ ) applied to the local homeomorphisms.

Important to our development will be the geometric logic of sets, i.e. that fragment of the internal intuitionistic logic of categories of sheaves that is preserved by inverse image functors. See [Vickers 99] for an account of how that includes not only the pure logic  $(\bigvee, \wedge, =, \exists)$  but also set constructions (principally colimits, finite limits, free algebras). We shall consistently understand *finite* to mean *Kuratowski finite*, and then the finite power set  $\mathcal{F}$  (free semilattice) and finitely bounded universal quantification are also geometric.

If  $T$  is a geometric theory, we write  $[T]$  for its classifying topos (as generalized space) and  $\mathcal{S}[T]$  for the corresponding category of sheaves (generalized universe of sets).

Very important to us will be the fact that locale theory, insofar as it is constructive, can be conducted internally in categories of sheaves  $\mathcal{S}B$  ( $B$  a topos). From an internal frame  $\Omega X$  in  $\mathcal{S}B$  one can construct a category of sheaves  $\mathcal{S}X$  with a *localic* geometric morphism  $f : X \rightarrow B$  (i.e. what [JoyTie 84] Section VI.5 call a *spatial* geometric morphism). The frame  $\Omega X$  can be recovered (up to isomorphism) as  $f_*(\Omega_X)$ , where  $\Omega_X$  is the subobject classifier in  $\mathcal{S}X$ . An important fact is that continuous maps between the locales (in the sense of frame homomorphisms between the frames) are equivalent to geometric morphisms between the toposes:

**2.1. THEOREM.** *Let  $f : X \rightarrow B$  be a map (geometric morphism) that is localic. Then for any other map  $g : E \rightarrow B$ , there is an equivalence between maps from  $E$  to  $X$  over  $B$ , and frame homomorphisms from  $f_*\Omega_X$  to  $g_*\Omega_E$  in  $\mathcal{S}B$ . ( $\Omega_X$  and  $\Omega_E$  are the subobject classifiers in  $\mathcal{S}X$  and  $\mathcal{S}E$ .)*

**PROOF.** Joyal and Tierney [JoyTie 84]. ■

**STAGES OF DEFINITION.** We now turn to the spatial aspects of locale theory. As is well-known, a locale  $X$  may not have an adequate supply of points in the sense of maps  $1 \rightarrow X$  (*global* points), and so reasoning solely in terms of those points cannot give full results about the locale. Instead, we must consider *generalized* points, maps  $Y \rightarrow X$ , where  $Y$ , an arbitrary topos, is the *stage of definition*. In particular, the identity map  $\text{Id} : X \rightarrow X$  can be considered a point (the *generic* point) in the internal logic of  $\mathcal{S}X$ , and for many purposes it is already a sufficiency of points. For instance, as discussed at length in [Vickers 99] and [Vickers 01], a map  $f : X \rightarrow Y$  can be described in full simply by giving a geometric construction of  $f(x)$  for the generic point  $x$  of  $X$ . The fact that this construction is intuitionistic means that it can be done internally in  $\mathcal{S}X$ , thus obtaining a point of  $Y$  there, and the fact that it is geometric means that it is preserved by pullback along geometric morphisms and so the same construction applies at any stage of definition.

Let us examine the relationship between global locale theory and locale theory at some stage  $B$ .

If  $X$  is a (global) locale, then we have a corresponding locale over  $B$ , namely the first projection, a localic map,  $B \times X \rightarrow B$ . This can be expressed concretely in frames, but in Section 5 we shall see how to express it using frame presentations.

Suppose  $Y$  is another global locale. A map from  $X$  to  $Y$  *at stage* (or *over*)  $B$  is a map from  $B \times X$  to  $B \times Y$  making this triangle commute:

$$\begin{array}{ccc} B \times X & \longrightarrow & B \times Y \\ & \searrow & \downarrow \\ & & B \end{array}$$

Globally, this is equivalent to a map from  $B \times X$  to  $Y$ . The correspondence respects order, for a pair of maps  $f \sqsubseteq g : X \rightarrow Y$  is equivalent to a map  $\mathbb{S} \times X \rightarrow Y$ . Working over  $B$ , this is equivalent to  $(B \times \mathbb{S}) \times_B (B \times X) \cong B \times \mathbb{S} \times X \rightarrow B \times Y$  over  $B$  (it follows from the main result of Section 5 that the Sierpinski locale as calculated at stage  $B$  is



just  $B \times \mathbb{S} \rightarrow B$ ), and that is equivalent to  $\mathbb{S} \times B \times X \rightarrow Y$ , i.e. two maps  $B \times X \rightarrow Y$  with one less than the other.

The final locale at stage  $B$  is the identity map  $B \cong B \times 1 \rightarrow B$ , and so a point of  $X$  at stage  $B$ , a map  $1 \rightarrow X$  at stage  $B$ , is, globally, a map  $B \rightarrow X$ .

An open of  $X$  at stage  $B$  is a map  $X \rightarrow \mathbb{S}$  at stage  $B$  and hence, globally, a map  $B \times X \rightarrow \mathbb{S}$ , in other words an open of  $B \times X$ .

If, at stage  $B$ , we have a point  $x$  and an open  $a$  of  $X$ , then we write  $x \vDash a$  for the truth value (point of  $\mathbb{S}$ ) calculated over  $B$  as  $x; a : 1 \rightarrow X \rightarrow \mathbb{S}$ . Globally, this is given as  $\langle B, x \rangle; a : B \rightarrow B \times X \rightarrow \mathbb{S}$ .

**2.2. PROPOSITION.** *Let  $X$  be a locale, and let  $a$  and  $b$  be two sublocales. Then  $a \leq b$  iff every point of  $a$  is also in  $b$ .*

**PROOF.** If the generic point of  $a$  is also in  $b$ , then the same goes for every point at every stage. Working with the generic point is equivalent to working over  $a$ . Viewed globally, the generic point  $\Delta : a \rightarrow a \times a$  is, as point of  $X$ ,  $\langle \text{Id}_a, i_a \rangle : a \rightarrow a \times X$  where  $i_a : a \hookrightarrow X$  is the sublocale inclusion. For this to be in  $b$  we must have some  $\langle \text{Id}_a, f \rangle : a \rightarrow a \times b$  with  $\langle \text{Id}_a, f \rangle; \langle \text{Id}_a, i_b \rangle = \langle \text{Id}_a, i_a \rangle$ . In other words  $i_a = f; i_b$ , i.e.  $a \leq b$ . ■

**SUPLATTICES AND PREFRAMES.** A *suplattice* [JoyTie 84] is a complete join semilattice; the homomorphisms are the functions that preserve all joins. A *preframe* [JoVic 91] is a poset equipped with finite meets and directed joins, the former distributing over the latter. The homomorphisms are the functions that preserve finite meets and directed joins. In various contexts we shall use “SupL” and “PreFr” as abbreviations for these.

**CONTINUOUS DCPOS.** We shall need some results on continuous dcpos (or continuous posets), treated as locales. To access geometric methods for handling them, we present them using the continuous information systems of [Vickers 93].

**2.3. DEFINITION.** *A continuous information system is a set  $P$  equipped with an idempotent binary relation  $<$ .*

“Idempotent” means that  $<; < = <$ . The  $\subseteq$  direction of this says that  $<$  is transitive, while the  $\supseteq$  direction says that it is *interpolative* — if  $x < z$  then  $x < y < z$  for some  $y$ .

**2.4. DEFINITION.** *Let  $P$  be a continuous information system. Then an ideal of  $P$  is a subset  $I$  such that*

1.  $I$  is inhabited.
2.  $I$  is lower closed under  $<$ .
3. If  $x, y \in I$  then there is some  $z \in I$  such that  $x < z$  and  $y < z$ .

(Note from (3) that if  $x \in I$  then there is some  $z \in I$  such that  $x < z$ : in other words,  $I$  is rounded.)

The ideal completion  $\text{Idl}(P)$  of  $P$  is the locale whose points are the ideals of  $P$ . It can be presented as

$$\begin{aligned}\Omega \text{Idl}(P) &= \text{Fr}\langle \uparrow x \ (x \in P) \mid \mathbf{true} \leq \bigvee_{x \in P} \uparrow x \\ &\quad \uparrow y \leq \uparrow x \ (x < y) \\ &\quad \uparrow x \wedge \uparrow y \leq \bigvee \{ \uparrow z \mid x < z, y < z \} \rangle\end{aligned}$$

where for any ideal  $I$  we have  $I \models \uparrow x$  iff  $x \in I$ .

### 3. The upper and lower powerlocales

The powerlocales  $P_L$  (lower) and  $P_U$  (upper) are conveniently summarized in [Vickers 97]. They are the functor parts of two monads on  $\mathbf{Loc}$ , denoted  $(P_L, \downarrow, \sqcup)$  and  $(P_U, \uparrow, \sqcap)$ .

3.1. DEFINITION. *The lower powerlocale is defined by*

$$\begin{aligned}\Omega P_L X &= \text{Fr}\langle \Omega X \text{ (qua SupL)} \rangle \\ &= \text{Fr}\langle \diamond a \ (a \in \Omega X) \mid \diamond \text{ preserves all joins} \rangle \\ \Omega P_L(f)(\diamond b) &= \diamond \Omega f(b) \ (f : X \rightarrow Y) \\ \Omega(\downarrow)(\diamond a) &= a \\ \Omega(\sqcup)(\diamond a) &= \diamond \diamond a\end{aligned}$$

(Its points are the weakly closed sublocales of  $X$  with open domain.)

3.2. DEFINITION. *The upper powerlocale is defined by*

$$\begin{aligned}\Omega P_U X &= \text{Fr}\langle \Omega X \text{ (qua PreFr)} \rangle \\ &= \text{Fr}\langle \square a \ (a \in \Omega X) \mid \square \text{ preserves finite meets and directed joins} \rangle \\ \Omega P_U(f)(\square b) &= \square \Omega f(b) \ (f : X \rightarrow Y) \\ \Omega(\uparrow)(\square a) &= a \\ \Omega(\sqcap)(\square a) &= \square \square a\end{aligned}$$

(Its points are the compact fitted sublocales of  $X$ .)

Note a crucial feature of these definitions: they are *not* geometric in form. This is because neither suplattice structure nor preframe structure is preserved by inverse image functors. A substantial part of this paper is devoted to showing that, nonetheless, the powerlocales *are* geometric.

An important feature is that  $P_L$  is a *KZ-monad* and  $P_U$  is a *coKZ-monad* (i.e. a KZ-monad on  $\mathbf{Loc}^{co}$ , with the order reversed), so we first mention some of the standard properties of these.

3.3. DEFINITION. *Let  $\mathcal{C}$  be a poset-enriched category, and let  $(T, \eta, \mu)$  be a monad on  $\mathcal{C}$  for which  $T$  preserves the order on morphisms. Then the monad is a KZ-monad iff, for each  $X$ ,  $T\eta_X \sqsubseteq \eta_{TX}$ .*

3.4. PROPOSITION. *Let  $\mathcal{C}$  be a poset-enriched category, and let  $(T, \eta, \mu)$  be a KZ-monad on  $\mathcal{C}$ .*

1. *For each object  $X$ ,  $\mu_X$  is right adjoint to  $T\eta_X$  and left adjoint to  $\eta_{TX}$  – this amounts to two inequalities*

$$\mu; T\eta_X \sqsubseteq \text{Id}_{T^2X} \sqsubseteq \mu; \eta_{TX}$$

*(In fact, these conditions are also sufficient for the KZ-property.)*

2. *If  $X$  is an object and  $\alpha : TX \rightarrow X$  is  $T$ -algebra structure on  $X$ , then  $\alpha$  is left adjoint to  $\eta_X$ . Hence  $T$ -algebra structure on  $X$ , if it exists at all, is unique.*

PROOF. [Kock 95] ■

Dually, a monad  $(T, \eta, \mu)$  is *coKZ* iff  $T\eta_X \supseteq \eta_{TX}$ .

3.5. PROPOSITION.  $P_L$  is a KZ-monad.

PROOF. We omit the proof that it is a monad. To show it is KZ, we show

$$P_L(\downarrow_X) \sqsubseteq \downarrow_{P_L X}: P_L X \rightarrow P_L^2 X$$

The finite meets  $\bigwedge_{a \in S} \diamond a$  form a basis for  $P_L X$ , and so the opens  $\diamond \bigwedge_{a \in S} \diamond a$  form a subbasis of  $P_L^2 X$ . Examining inverse image functions, we find that the required inequality follows from

$$\diamond \bigwedge_{a \in S} a \leq \bigwedge_{a \in S} \diamond a$$

■

3.6. PROPOSITION.  $P_U$  is a coKZ-monad.

PROOF. To show it is coKZ, we show

$$P_U(\uparrow_X) \supseteq \uparrow_{P_U X}: P_U X \rightarrow P_U^2 X.$$

The finite joins  $\bigvee_{a \in S} \square a$  form a *preframe basis* for  $P_U X$ , in the sense that every open is a directed join of finite meets of them. This is because the  $\square b$ s form a subbasis, so every open is a join of finite meets of them. An arbitrary join is a directed join of finite joins, and by finite distributivity a finite join of finite meets is a finite meet of finite joins. Hence the opens  $\square \bigvee_{a \in S} \square a$  form a subbasis of  $P_U^2 X$ . Examining inverse image functions, we find that the required inequality follows from

$$\square \bigvee_{a \in S} a \geq \bigvee_{a \in S} \square a$$

■

3.7. PROPOSITION. *Let  $X$  be a locale. Then the following are equivalent.*

1. *The unique map  $! : X \rightarrow 1$  and the diagonal map  $\Delta : X \rightarrow X^2$  have right adjoints.*
2. *For every  $n$ , the diagonal  $\Delta : X \rightarrow X^n$  has a right adjoint.*
3.  *$X$  is a localic semilattice whose semilattice operations are meets with respect to the specialization order.*

PROOF.  $2 \Rightarrow 1$  a fortiori.

$1 \Rightarrow 3$ : Let us write  $\top$  and  $\sqcap$  for the right adjoints of  $!$  and  $\Delta$ . It is easy to show that these operations satisfy the semilattice laws. For instance, for associativity one shows that both sides of the equation give right adjoints for the diagonal  $\Delta : X \rightarrow X^3$ . To define  $n$ -ary meets with respect to the specialization order is to define a right adjoint  $\sqcap$  for the diagonal  $\Delta : X \rightarrow X^n$ : for we want  $\sqcap_{i=1}^n x_i \sqsubseteq y$  iff  $(x_i) \sqsubseteq \Delta(y)$ .

$3 \Rightarrow 2$  also now follows. ■

3.8. DEFINITION. *A locale  $X$  is a localic meet semilattice if it satisfies the equivalent conditions of Proposition 3.7.*

Proposition 3.7 dualizes, replacing right adjoints by left adjoints and meets by joins. A locale satisfying the dual proposition is a *localic join semilattice*.

3.9. PROPOSITION.

1. *Any  $P_U$ -algebra is a localic meet semilattice.*
2. *Any  $P_L$ -algebra is a localic join semilattice.*

PROOF. 1. Let  $X$  be a locale. We first show that  $P_U X$  is a localic meet semilattice. We define  $\top : 1 \rightarrow P_U X$  and  $\sqcap : (P_U X)^2 \rightarrow P_U X$  by

$$\begin{aligned}\Omega \top(\square a) &= \mathbf{true} \\ \Omega \sqcap(\square a) &= \square a \times \square a\end{aligned}$$

To show that  $\top$  is right adjoint to  $!$ , we calculate  $\top; ! \sqsubseteq \text{Id}_1$  (actually, equality is obvious) and  $\text{Id}_{P_U X} \sqsubseteq !; \top$ . The latter follows because

$$\square a \leq \mathbf{true} = \Omega!(\mathbf{true}) = \Omega! \circ \Omega \top(\square a).$$

For  $\sqcap$  right adjoint to  $\Delta$ , we want  $\sqcap; \Delta \sqsubseteq \text{Id}_{(P_U X)^2}$  and  $\text{Id}_{P_U X} \sqsubseteq \Delta; \sqcap$ . For the former, we note that the opens  $\square a$  of  $P_U X$  form a basis, and so the opens  $\square a \times \square b$  of  $(P_U X)^2$  form a basis. We then calculate

$$\begin{aligned}\Omega(\sqcap; \Delta)(\square a \times \square b) &= \Omega \sqcap(\square a \wedge \square b) = \Omega \sqcap(\square(a \wedge b)) \\ &= \square(a \wedge b) \times \square(a \wedge b) \leq \square a \times \square b\end{aligned}$$

For the latter,

$$\Omega(\Delta; \sqcap)(\sqcap a) = \Omega\Delta(\sqcap a \times \sqcap a) = \sqcap a \wedge \sqcap a = \sqcap a$$

Now let  $X$  be any  $P_U$ -algebra with structure map  $\alpha : P_U X \rightarrow X$ . We can define its semilattice operations  $\sqcap_n : X^n \rightarrow X$  by the composite

$$X^n \xrightarrow{\uparrow^n} (P_U X)^n \xrightarrow{\sqcap_n} P_U X \xrightarrow{\alpha} X$$

Now

$$\begin{aligned} \Delta; \uparrow^n; \sqcap_n; \alpha &= \uparrow; \Delta; \sqcap_n; \alpha = \uparrow; \alpha = \text{Id}_X \\ \uparrow^n; \sqcap_n; \alpha; \Delta &= \uparrow^n; \sqcap_n; \alpha; \Delta; \uparrow^n; \alpha^n \\ &= \uparrow^n; \sqcap_n; \alpha; \uparrow; \Delta; \alpha^n \\ &\sqsubseteq \uparrow^n; \sqcap_n; \Delta; \alpha^n \\ &\sqsubseteq \uparrow^n; \alpha^n = \text{Id}_{X^n} \end{aligned}$$

so that these are right adjoint to the diagonals as required.

2. The proof is more or less dual. The localic join semilattice operations  $\perp$  and  $\sqcup$  on  $P_L X$  are defined by

$$\begin{aligned} \Omega\perp(\diamond a) &= \mathbf{false} \\ \Omega\sqcup(\diamond a) &= \diamond a \odot \diamond a \end{aligned}$$

(In a product  $X \times Y$ , we write  $a \odot b$  for  $a \times Y \vee X \times b$ .) ■

We shall need results on powerlocales of dcpo's.

**3.10. DEFINITION.** *Let  $(X, <)$  be a continuous information system. Then the orderings  $<_U$  and  $<_L$  on  $\mathcal{F}G$ , the upper and lower orders, are defined by*

$$\begin{aligned} S <_U T &\text{ iff } \forall t \in T. \exists s \in S. s < t \\ S <_L T &\text{ iff } \forall s \in S. \exists t \in T. s < t \end{aligned}$$

*Both are idempotent. If in addition  $<$  is reflexive (so it is a preorder), then so are  $<_U$  and  $<_L$ , and we write  $\equiv_U$  and  $\equiv_L$  for the corresponding equivalence relations.*

**3.11. THEOREM.** *Let  $(X, <)$  be a continuous information system. Then –*

$$\begin{aligned} P_U(\text{Idl}(X)) &\cong \text{Idl}(\mathcal{F}X, <_U) \\ P_L(\text{Idl}(X)) &\cong \text{Idl}(\mathcal{F}X, <_L) \end{aligned}$$

**PROOF.** [Vickers 93]. ■

### 4. The double powerlocale $\mathbb{P}$

In [JoVic 91] it is shown that the upper and lower powerlocales commute up to isomorphism. We shall write  $\mathbb{P}$  for (without prejudice)  $P_U P_L$  or  $P_L P_U$ . Under the isomorphisms, the open  $\square \diamond a$  of  $P_U P_L X$  ( $a$  an open of  $X$ ) corresponds to the open  $\diamond \square a$  of  $P_L P_U X$ . We write  $\boxtimes a$  for the corresponding open of  $\mathbb{P}X$ ,  $\boxtimes$  being intended as a kind of combination of  $\square$  and  $\diamond$ .

Another way to understand the monad structure on  $P_U P_L$  is in terms of *distributivities* [Beck 69]:  $P_U$  and  $P_L$  distribute over each other.

4.1. DEFINITION. *Let  $(T_i, \eta_i, \mu_i)$  be a monad ( $i = 1, 2$ ). A distributive law of  $T_2$  over  $T_1$  is a natural transformation  $\sigma : T_2 \circ T_1 \rightarrow T_1 \circ T_2$  such that the following diagrams commute.*

$$\begin{array}{ccccc}
 T_2 & \xrightarrow{T_2 \eta_1} & T_2 \circ T_1 & \xleftarrow{\eta_2 T_1} & T_1 \\
 & \searrow & \downarrow \sigma & \swarrow & \\
 & \eta_1 T_2 & T_1 \circ T_2 & & \\
 \\
 T_2 \circ T_1 \circ T_1 & \xrightarrow{T_2 \mu_1} & T_2 \circ T_1 & \xleftarrow{\mu_2 T_1} & T_2 \circ T_2 \circ T_1 \\
 \downarrow \sigma T_1 & & & & \downarrow T_2 \sigma \\
 T_1 \circ T_2 \circ T_1 & & \downarrow \sigma & & T_2 \circ T_1 \circ T_2 \\
 \downarrow T_1 \sigma & & & & \downarrow \sigma T_2 \\
 T_1 \circ T_1 \circ T_2 & \xrightarrow{\mu_1 T_2} & T_1 \circ T_2 & \xleftarrow{T_1 \mu_2} & T_1 \circ T_2 \circ T_2
 \end{array}$$

There is then a monad structure on  $T_1 \circ T_2$ , with unit

$$\eta = \eta_1 \eta_2 : \text{Id} \rightarrow T_1 \circ T_2$$

and multiplication

$$\mu = (T_1 \sigma T_2); (\mu_1 \mu_2) : T_1 \circ T_2 \circ T_1 \circ T_2 \rightarrow T_1 \circ T_1 \circ T_2 \circ T_2 \rightarrow T_1 \circ T_2$$

4.2. PROPOSITION. *The isomorphism  $P_L P_U \cong P_U P_L$  is a distributive law of  $P_L$  over  $P_U$ . Its inverse is a distributive law of  $P_U$  over  $P_L$ .*

PROOF. The conditions are routine to check. ■

4.3. COROLLARY.  *$\mathbb{P}$  is the functor part of a monad on **Loc**, with unit  $\uparrow : X \rightarrow \mathbb{P}X$  defined by  $\Omega \uparrow (\boxtimes a) = a$  and multiplication  $\mathbf{H} : \mathbb{P}\mathbb{P}X \rightarrow \mathbb{P}X$  defined by  $\Omega \mathbf{H}(\boxtimes a) = \boxtimes \boxtimes a$ .*

(The notation is explained as follows. The powerlocales give monads  $(P_U, \uparrow, \sqcap)$  and  $(P_L, \downarrow, \sqcup)$ , where the notation is chosen to match the fact that the points of the powerlocales are sublocales of the original locale. Then, for instance, if  $x$  is a point of  $X$ , then the point  $\uparrow x$  of  $P_U X$  is the compact saturated sublocale of  $X$  whose points are those of  $X$  that are bigger than  $x$  in the specialization order. Then  $\uparrow$  is chosen as a combination of  $\uparrow$  and  $\downarrow$ , and  $\mathbf{H}$  as a combination of  $\sqcap$  and  $\sqcup$ .)

4.4. THEOREM.  $\Omega \mathbb{P}X \cong \text{Fr}\langle \Omega X \text{ (qua dcpo)} \rangle$ .

PROOF. [JoVic 91]. ■

Our contention is that the  $\mathbb{P}$ -algebras are in some sense “localic frames”. They are both  $P_U$ -algebras and  $P_L$ -algebras, and this makes them both localic meet-semilattices and localic join semilattices. Moreover, as lattices they are distributive. But every locale has directed joins (of its points), and the meets, being continuous, preserve those directed joins. Hence we can reasonably think of a  $\mathbb{P}$ -algebra as having all joins as well as frame distributivity.

We shall show in due course that the Sierpinski locale  $\mathbb{S}$ , whose points are the subsets of  $1$ , is the initial  $\mathbb{P}$ -algebra (in fact,  $\mathbb{S} \cong \mathbb{P}\emptyset$ ). Also, for locale  $X$ , the exponential  $\mathbb{S}^X$  is a  $\mathbb{P}$ -algebra if it exists. Then the points of  $X$  are equivalent to the  $\mathbb{P}$ -algebra homomorphisms from  $\mathbb{S}^X$  to  $\mathbb{S}$ , and we have lifted to **Loc** the standard fact that the points of  $X$  are equivalent to the frame homomorphisms from  $\Omega X$  to  $\Omega$ .

In a frame the meets and joins are determined by the order. In a locale the specialization order is intrinsic, and so we should expect that  $\mathbb{P}$ -algebra structure, where it exists at all, is unique. We shall now prove this.

4.5. PROPOSITION. *Let  $\alpha : \mathbb{P}X \rightarrow X$  make  $X$  a  $\mathbb{P}$ -algebra.*

1.  *$X$  is a  $P_U$ -algebra and  $\alpha$  is a  $P_U$ -homomorphism.*
2.  *$X$  is a  $P_L$ -algebra and  $\alpha$  is a  $P_L$ -homomorphism.*
3.  *$\alpha$  is the unique  $\mathbb{P}$ -algebra structure of  $X$ .*

PROOF. 1. Define  $\alpha_U = P_U(\downarrow); \alpha : P_U X \rightarrow P_U P_L X = \mathbb{P}X \rightarrow X$ . The properties of distributivities tell us that

$$\uparrow = (\uparrow; P_U(\downarrow)) : X \rightarrow P_U X \rightarrow P_U P_L X = \mathbb{P}X,$$

and it follows that

$$\uparrow; \alpha_U = \uparrow; \alpha = \text{Id}_X$$

Next, to complete the proof that  $\alpha_U$  makes  $X$  a  $P_U$ -algebra, we should prove that  $\square; \alpha_U = P_U(\alpha_U); \alpha_U$ . However, it is easier to prove first that  $\alpha$  is a  $P_U$ -homomorphism (note that  $\mathbb{P}X = P_U P_L X$  is itself a  $P_U$ -algebra), i.e. that  $\square; \alpha = P_U \alpha; \alpha_U$ . For this, note first that

$$\square = P_U(\downarrow); \mathbf{H} : P_U \mathbb{P}X \rightarrow \mathbb{P}^2 X \rightarrow \mathbb{P}X$$

for both have inverse image map  $\boxtimes a \mapsto \square \square \diamond a$ . It follows that

$$\begin{aligned} \square; \alpha &= P_U(\downarrow); \mathbf{H}; \alpha \\ &= P_U(\downarrow); \mathbb{P}\alpha; \alpha \\ &= P_U(\downarrow; P_L \alpha); \alpha \\ &= P_U(\alpha; \downarrow); \alpha \\ &= P_U \alpha; P_U(\downarrow); \alpha \\ &= P_U \alpha; \alpha_U \end{aligned}$$

Now to complete the proof that  $\alpha_U$  makes  $X$  a  $P_U$ -algebra, we have

$$\begin{aligned} \sqcap; \alpha_U &= \sqcap; P_U(\downarrow); \alpha \\ &= P_U^2(\downarrow); \sqcap; \alpha \\ &= P_U^2(\downarrow); P_U \alpha; \alpha_U \\ &= P_U(P_U(\downarrow); \alpha); \alpha_U \\ &= P_U(\alpha_U); \alpha_U \end{aligned}$$

The proof of (2) is formally very similar, using  $\alpha_L = P_L(\uparrow); \alpha : P_L X \rightarrow P_L P_U X \cong \mathbb{P}X \rightarrow X$ .

For (3), suppose  $\beta : \mathbb{P}X \rightarrow X$  also makes  $X$  a  $\mathbb{P}$ -algebra. By Proposition 3.4 we have  $\alpha_U = \beta_U$  and  $\alpha_L = \beta_L$ . Let us fix on  $\mathbb{P}X$  as  $P_U P_L X$  and be careful about the isomorphism with  $P_L P_U X$ .  $\alpha_L$  is strictly speaking the map

$$\alpha_L = P_L(\uparrow); \cong; \alpha : P_L X \rightarrow P_L P_U X \cong P_U P_L X = \mathbb{P}X \rightarrow X$$

We can see that  $P_L(\uparrow); \cong$  is equal to  $\uparrow$ , since the inverse image functions for both take  $\square \diamond a \mapsto \diamond a$ . It follows that  $\alpha$  is the unique  $P_U$ -homomorphism from  $P_U P_L X$  to  $X$  extending  $\alpha_L$ . Similarly  $\beta$  is the unique  $P_U$ -homomorphism from  $P_U P_L X$  to  $X$  extending  $\beta_L$ , and since  $\alpha_L = \beta_L$  we deduce  $\alpha = \beta$ . ■

The following result is analogous to the fact that any order isomorphism between frames is a frame isomorphism.

4.6. COROLLARY. *Any homeomorphism between  $\mathbb{P}$ -algebras is a  $\mathbb{P}$ -algebra isomorphism.*

PROOF. Let  $\alpha : \mathbb{P}X \rightarrow X$  and  $\beta : \mathbb{P}Y \rightarrow Y$  be  $\mathbb{P}$ -algebras, and let  $f : X \rightarrow Y$  be a homeomorphism. It is straightforward to verify that  $\mathbb{P}f; \beta; f^{-1}$  is a  $\mathbb{P}$ -algebra structure on  $X$ , so  $\mathbb{P}f; \beta; f^{-1} = \alpha$ . ■

4.7. PROPOSITION. *The initial  $\mathbb{P}$ -algebra  $\mathbb{P}\emptyset$  is  $\mathbb{S}$ . Its structure map, from  $\mathbb{P}\mathbb{S}$  to  $\mathbb{S}$ , corresponds to the open  $\boxtimes\{\top\}$ .*

PROOF. This is most conveniently calculated applying Theorem 3.11 in the special case of algebraic dcpos,  $\text{Idl } P$  where  $P$  is a poset.

Now the initial locale  $\emptyset$  is the ideal completion of the empty set, so  $P_L \emptyset$  is the ideal completion of  $\{\emptyset\}$ . This is just the final, one-point locale. Then  $P_U P_L \emptyset$  is the ideal completion of a 2-element poset  $\{\perp \sqsubseteq \top\}$ . A point of  $P_U P_L \emptyset$  is an ideal of  $\{\perp, \top\}$ , i.e. a subset  $I \subseteq \{\perp, \top\}$  that is directed and lower closed. Since it is inhabited and lower closed it must contain  $\perp$ , so it is completely determined by whether it contains  $\top$  — in other words, by its intersection with  $\{\top\}$ . But if  $I'$  is a subset of  $\{\top\}$  then  $\{\perp\} \cup I'$  is an ideal. Hence there is a bijection between points of  $P_U P_L \emptyset$  and subsets of 1. Moreover, this bijection is geometrically definable. The Sierpinski locale  $\mathbb{S}$  is defined by its frame being free on one generator, and it follows that its points are the subsets of 1. We therefore have that  $\mathbb{P}\emptyset$  is homeomorphic to  $\mathbb{S}$ .



The generator  $\{\top\}$  of  $\Omega\mathcal{S}$  is  $\square\mathbf{false}$  in  $\mathbb{P}_U 1$ , i.e.  $\square\diamond\mathbf{false}$  in  $\mathbb{P}_U \mathbb{P}_L \emptyset$ , i.e.  $\boxtimes\mathbf{false}$  in  $\mathbb{P}\emptyset$ . The structure map is just  $\mathbf{H}$ , and the corresponding open is  $\Omega\mathbf{H}(\{\top\}) = \Omega\mathbf{H}(\boxtimes\mathbf{false}) = \boxtimes\boxtimes\mathbf{false} = \boxtimes\{\top\}$ . ■

We now investigate the localic lattice structure of  $\mathbb{P}$ -algebras.

**4.8. PROPOSITION.** *Let  $X$  be a  $\mathbb{P}$ -algebra. Then  $X$  is both a localic meet semilattice and a localic join semilattice, and the resulting lattice structure is distributive.*

**PROOF.** We have already shown that  $X$  is both a  $\mathbb{P}_U$ -algebra and a  $\mathbb{P}_L$ -algebra. We show distributivity first for algebras  $\mathbb{P}Y$ . We must consider two maps  $(\mathbb{P}Y)^3 \rightarrow \mathbb{P}Y$ . The first takes  $(u, v, w)$  to  $u \sqcap (v \sqcup w)$ , and its inverse image function takes

$$\boxtimes a \mapsto \boxtimes a \times \boxtimes a \mapsto \boxtimes a \times (\boxtimes a \odot \boxtimes a)$$

The second takes  $(u, v, w)$  to  $(u \sqcap v) \sqcup (u \sqcap w)$ . Its inverse image function takes

$$\begin{aligned} \boxtimes a &\mapsto \boxtimes a \odot \boxtimes a \mapsto (\boxtimes a \times \boxtimes a) \odot (\boxtimes a \times \boxtimes a) \\ &= ((\boxtimes a \times \mathbb{P}Y \times \mathbb{P}Y \times \mathbb{P}Y) \wedge (\mathbb{P}Y \times \boxtimes a \times \mathbb{P}Y \times \mathbb{P}Y)) \\ &\vee ((\mathbb{P}Y \times \mathbb{P}Y \times \boxtimes a \times \mathbb{P}Y) \wedge (\mathbb{P}Y \times \mathbb{P}Y \times \mathbb{P}Y \times \boxtimes a)) \\ &\mapsto ((\boxtimes a \times \mathbb{P}Y \times \mathbb{P}Y) \wedge (\mathbb{P}Y \times \boxtimes a \times \mathbb{P}Y)) \vee ((\boxtimes a \times \mathbb{P}Y \times \mathbb{P}Y) \wedge (\mathbb{P}Y \times \mathbb{P}Y \times \boxtimes a)) \\ &= (\boxtimes a \times \mathbb{P}Y \times \mathbb{P}Y) \wedge ((\mathbb{P}Y \times \boxtimes a \times \mathbb{P}Y) \vee (\mathbb{P}Y \times \mathbb{P}Y \times \boxtimes a)) \\ &= \boxtimes a \times (\boxtimes a \odot \boxtimes a) \end{aligned}$$

Hence the two are equal, and  $\mathbb{P}Y$  is distributive. Since the corresponding maps for  $X$  are got by composing the ones for  $\mathbb{P}X$  with  $\uparrow^3: X^3 \rightarrow (\mathbb{P}X)^3$  and  $\alpha: \mathbb{P}X \rightarrow X$ , it follows that the arbitrary  $\mathbb{P}$ -algebra  $X$  is also distributive. ■

Since  $\sqcap: X^2 \rightarrow X$  is continuous, it also distributes over directed joins in  $X$ , and this justifies our claim that  $\mathbb{P}$ -algebras are localic frames.

## 5. Geometricity of presentations

With this section we begin a series of geometricity results, essential for the geometric reasoning in the main results about  $\mathbb{P}$  and exponentiation. We start by discussing the geometricity of *frame presentations*. This is important, because frames themselves are not geometric.

A frame presentation will include sets  $G$  and  $R$  of generators and relations, and each relation can be written in the form  $e_1 \leq e_2$ , where  $e_1$  and  $e_2$  are frame expressions in the generators. Using frame distributivity, each  $e_i$  can be written as a join of finite meets of generators; and then the relation can be replaced by a set of relations, one for each disjunct in  $e_1$ , saying that the disjunct is  $\leq e_2$ . After all this rewriting we have that each relation  $r$  is of the form –

$$\text{finite meet of generators} \leq \text{join of finite meets of generators}$$

Let us write  $\lambda(r) \in \mathcal{F}G$  for the finite set of conjuncts on the left. For the right-hand side, we have an arbitrary set of disjuncts: so what we should do is take the set  $D$  of all disjuncts in all relations, fibred over  $R$  by some  $\pi : D \rightarrow R$ . Each disjunct  $d$  is a conjunction of a finite set  $\rho(d)$  of generators, so the relation  $r$  has been formalized as –

$$\bigwedge \lambda(r) \leq \bigvee_{\pi(d)=r} \bigwedge \rho(d)$$

5.1. DEFINITION. A frame presentation (or a  $(G,R,D)$ -system) is a structure comprising three sets  $G$ ,  $R$  and  $D$  with functions  $\lambda : R \rightarrow \mathcal{F}G$ ,  $\rho : D \rightarrow \mathcal{F}G$ , and  $\pi : D \rightarrow R$ .

We write  $\text{FrPr}$  for the (geometric) theory of frame presentations.

Given such a frame presentation, then in conformity with established notation we shall write  $\text{Fr}\langle G|R \rangle$  for the frame presented by it. (Of course, this conceals the involvement of  $D$ ,  $\lambda$ ,  $\rho$  and  $\pi$ .) Then a frame homomorphism from  $\text{Fr}\langle G|R \rangle$  to a frame  $A$  is given by a function  $\gamma : G \rightarrow A$  that respects the relations in the following way. Since  $A$  is a semilattice under  $\wedge$ , and  $\mathcal{F}G$  is the free semilattice over  $G$ ,  $\gamma$  extends uniquely to a semilattice homomorphism  $\gamma' : (\mathcal{F}G, \cup) \rightarrow (A, \wedge)$  such that  $\gamma'(\{g\}) = \gamma(g)$ . We want for each relation  $r$  that

$$\gamma' \circ \lambda(r) \leq \bigvee \{ \gamma' \circ \rho(d) : \pi(d) = r \}$$

5.2. THEOREM. Let  $T$  be a geometric theory whose ingredients include a frame presentation as above, and let  $A$  be the frame in  $\mathcal{S}[T]$  presented by it. Then the corresponding locale over  $[T]$  classifies the theory  $T'$  that is  $T$  extended by –

- a predicate symbol  $I(g)$  ( $g : G$ )
- an axiom

$$\forall r : R. ((\forall g \in \lambda(r). I(g)) \longrightarrow \exists d : D. (\pi(d) = r \wedge \forall g \in \rho(d). I(g)))$$

PROOF. Let  $f : E \rightarrow [T]$  be any topos over  $[T]$ . We know (by Theorem 2.1) that geometric morphisms over  $[T]$  from  $E$  to the locale are equivalent to frame homomorphisms in  $\mathcal{S}[T]$  from  $A$  to  $f_*\Omega_E$ , and these are equivalent to functions from  $G$  to  $f_*\Omega_E$  that respect the relations. On the other hand, geometric morphisms over  $[T]$  from  $E$  to  $[T']$  are equivalent to subsets of  $f^*(G)$  that satisfy the axiom. We show that these are equivalent, and that suffices to show that the locale and  $[T']$  are equivalent over  $[T]$ .

Functions  $\gamma : G \rightarrow f_*\Omega_E$  are equivalent to functions from  $f^*G$  to  $\Omega_E$ , which in turn are equivalent to subsets  $I$  of  $f^*G$ . The difficult part is to show that the one respects the relations iff the other satisfies the axiom.

Recall some general properties about how  $f^*$  and  $f_*$  relate to algebras for any finitary algebraic theory. First, because both  $f_*$  and  $f^*$  preserve finite products, they transform algebras into algebras. (In particular, this gives the distributive lattice structure on  $f_*\Omega_E$ , though not the frame structure – for Mikkelsen’s description of joins in  $f_*\Omega_E$  see [Johnstone 77], Proposition 5.36.) Moreover, if  $X$  and  $Y$  are algebras in  $\mathcal{S}[T]$  and  $\mathcal{S}E$

respectively, and  $\theta : X \rightarrow f_*Y$  and  $\phi : f^*X \rightarrow Y$  are adjoint transposes of each other, then  $\theta$  is a homomorphism iff  $\phi$  is. Finally, if  $\mathbb{F}$  denotes the free algebra construction, then  $f^*$  preserves it:  $f^*\mathbb{F}X \cong \mathbb{F}(f^*X)$  by an algebra isomorphism in  $\mathcal{S}E$  (we shall apply this for the theory of semilattices, so that  $\mathbb{F}$  is  $\mathcal{F}$ ). We interrupt the proof of the theorem with a lemma about finitely bounded universal quantification, ■

**5.3. LEMMA.** *Let  $\theta : X \rightarrow f_*\Omega_E$  in  $\mathcal{S}[T]$  give semilattice homomorphism  $\theta' : \mathcal{F}X \rightarrow (f_*\Omega_E, \wedge)$ , and let their adjoint transposes be  $\phi : f^*X \rightarrow \Omega_E$  and  $\phi' : f^*\mathcal{F}X \rightarrow \Omega_E$ . Then the function*

$$\cong; \phi' : \mathcal{F}(f^*X) \cong f^*\mathcal{F}X \rightarrow \Omega_E$$

*maps  $S$  to the truth value  $[\forall x \in S. \phi(x)]$ .*

PROOF.

$$\begin{array}{ccc} f^*X & \xrightarrow{\phi} & \Omega_E \\ \mathcal{F}f^*X \swarrow \{ - \} & \downarrow f^*\{ - \} & \nearrow \phi' \\ & f^*\mathcal{F}X & \\ \cong \nearrow & & \\ & & \end{array} \quad \begin{array}{ccc} X & \xrightarrow{\theta} & f_*\Omega_E \\ \{ - \} \downarrow & \nearrow \theta' & \\ \mathcal{F}X & & \end{array}$$

In the left-hand diagram, both triangles commute: the left-hand one by definition of the isomorphism, and the right-hand one by naturality of the adjoint transpose. Also,  $\phi'$  is a semilattice homomorphism ( $\Omega_E$  as  $\wedge$ -semilattice) because  $\theta'$  is. It follows that  $\cong; \phi'$  is the unique semilattice homomorphism mapping each  $\{x\}$  to  $\phi(x)$ , but  $S \mapsto [\forall x \in S. \phi(x)]$  is such a one. ■

PROOF. We now return to the proof of the theorem. From the lemma (applied to  $\gamma : G \rightarrow f_*\Omega_E$ ), and using naturality, it follows that the adjoint transposes of  $\lambda; \gamma' : R \rightarrow f_*\Omega_E$  and  $\rho; \gamma' : D \rightarrow f_*\Omega_E$  correspond to the subsets  $\{r : \forall g \in f^*\lambda(r). I(g)\}$  and  $\{d : \forall g \in f^*\rho(d). I(g)\}$  of  $f^*R$  and  $f^*D$ . It remains to show that the adjoint transpose of the function  $r \mapsto \bigvee \{ \gamma' \circ \rho(d) : \pi(d) = r \}$  from  $R$  to  $f_*\Omega_E$  corresponds to the subset

$$\{r : \exists d. (f^*\pi(d) = r \wedge \forall g \in f^*\rho(d). I(g))\}$$

of  $f^*R$ . Our function from  $R$  to  $f_*\Omega_E$  calculates the join of a generalized element of  $\Omega_{[T]}^{f_*\Omega_E}$ , namely  $r \mapsto \{ \gamma' \circ \rho(d) : \pi(d) = r \}$ . The generalized element comes from a function from  $R \times f_*\Omega_E$  to  $\Omega_{[T]}$ , and this corresponds to a subset of  $R \times f_*\Omega_E$ , namely the image of  $\langle \pi, \gamma' \circ \rho \rangle : D \rightarrow R \times f_*\Omega_E$ . We obtain a function  $\langle f^*\pi, \alpha \rangle : f^*D \rightarrow f^*R \times \Omega_E$ , where  $\alpha : f^*D \rightarrow \Omega_E$  is the adjoint transpose of  $\gamma' \circ \rho$ , which by the lemma corresponds to the subset  $\{d : \forall g \in f^*\rho(d). I(g)\}$  of  $f^*D$ . The image of  $\langle f^*\pi, \alpha \rangle$  gives a function from  $f^*R$  to  $\Omega_E^{\Omega_E}$ ,

$$r \mapsto \{a \in \Omega_E \mid \exists d : f^*D. (f^*\pi(d) = r \wedge a = [\forall g \in f^*\rho(d). I(g)])\}$$

and the join of this set is the truth value  $[\exists d : f^*D. (f^*\pi(d) = r \wedge \forall g \in f^*\rho(d). I(g))]$ , and Mikkelsen's description says that this is the adjoint transpose of

$$r \mapsto \bigvee \{ \gamma' \circ \rho(d) \mid \pi(d) = r \}.$$

■

5.4. COROLLARY. *Let  $E$  be a topos, let  $P = (G, R, D, \lambda, \rho, \pi)$  be a frame presentation in  $\mathcal{S}E$  and let  $X \rightarrow E$  be the corresponding locale over  $E$ . Let  $f : E' \rightarrow E$  be a map. Then the locale  $X'$  over  $E'$  corresponding to the frame presentation  $f^*(P)$  in  $\mathcal{S}E'$  is the pullback  $E' \times_E X$ .*

PROOF. Let  $[\text{FrPr}][\text{pt}]$  classify a frame presentation together with a point described by predicate  $I$  and axiom as in Theorem 5.2. The map  $[\text{FrPr}][\text{pt}] \rightarrow [\text{FrPr}]$  that forgets the point is localic, and is in effect the generic presented locale. The theorem tells us that  $X \simeq E \times_{[\text{FrPr}]} [\text{FrPr}][\text{pt}]$  and  $X' \simeq E' \times_{[\text{FrPr}]} [\text{FrPr}][\text{pt}] \simeq E' \times_E X$ . ■

Corollary 5.4 enables us to address the geometricity of locale constructions in the following way. Suppose  $T : \mathbf{Loc} \rightarrow \mathbf{Loc}$  is a functorial construction of locales that can be described intuitionistically. Then for any topos  $E$  we have a functor  $T_E : \mathbf{Loc}/E \rightarrow \mathbf{Loc}/E$ . Now suppose we have a map  $p : E' \rightarrow E$ . If we have a locale  $X$  over  $E$ , then we have two locales over  $E'$ , namely  $p^*(T_E(X))$  and  $T_{E'}(p^*(X))$ . How do they relate to each other? In general there is no obvious morphism between them. However, we shall be interested in the case where they are isomorphic. In this case we shall say  $T$  is *geometric*, in that it is preserved by pullback along maps.

Corollary 5.4 allows us to use frame presentations to reduce the problem to geometric constructions on *sets*, in other words those set constructions that are preserved by the inverse image functors. As remarked earlier, the principal geometric set constructions are colimits, finite limits and free algebra constructions, and also universal quantification bounded over (Kuratowski) finite sets.

5.5. PROPOSITION. *Let  $T : \mathbf{Loc} \rightarrow \mathbf{Loc}$  be an intuitionistically describable functorial construction of locales as above, and let  $T'$  be a geometric construction on frame presentations (in other words, an endomorphism  $T' : [\text{FrPr}] \rightarrow [\text{FrPr}]$  of the classifying topos for frame presentations). Suppose we have an intuitionistic proof (valid over any topos) that if a  $(G, R, D)$ -system  $P$  presents a locale  $X$  then  $T'(P)$  presents (up to homeomorphism)  $T(X)$ . Then  $T$  is geometric.*

PROOF. We shall write  $T'_E$  and  $T_E$  for  $T'$  and  $T$  working at stage  $E$ . Let  $f : E' \rightarrow E$  be a map of toposes, and suppose, working over  $E$ , that  $X$  is a locale presented by  $P$ . Then by Corollary 5.4,  $f^*X$  is presented over  $E'$  by  $f^*(P)$ . Hence  $T_{E'}(f^*X)$  is presented by  $T'_{E'}(f^*P)$ . By geometricity of  $T'$  this is isomorphic to  $f^*(T'_E(P))$ , and that, using Corollary 5.4 again, presents  $f^*(T_E(X))$ . Hence  $f^*(T_E(X)) \cong T_{E'}(f^*X)$  as desired. ■

In what follows, the set of generators will often have structure (partial order, or meets or joins) that is to be preserved in the frame generated. This can easily be incorporated into the relations and disjuncts, and we use the “qua” notation to indicate that this should be done.

5.6. DEFINITION. *Let  $(G, R, D, \lambda, \rho, \pi)$  be a GRD-system in which  $G$  is equipped with a preorder  $\leq_G$ . Then by*

$$(G \text{ (qua preorder)}, R, D, \lambda, \rho, \pi)$$

we denote the GRD-system

$$(G, R', D', \lambda', \rho', \pi')$$

in which

$$\begin{aligned} R' &= R+ \leq_G \\ \lambda'(r) &= \lambda(r), \lambda'(g_1, g_2) = \{g_1\} \\ D' &= D+ \leq_G \\ \rho'(d) &= \rho(d), \rho'(g_1, g_2) = \{g_2\} \\ \pi' &= \pi+ \leq_G \end{aligned}$$

This presents

$$\begin{aligned} \text{Fr}\langle G | \bigwedge \lambda(r) \leq \bigvee_{\pi(d)=r} \bigwedge \rho(d) \ (r \in R), \\ g_1 \leq g_2 \ (g_1 \leq g_2 \text{ in } G) \rangle \end{aligned}$$

which is isomorphic to  $\text{Fr}\langle G \text{ (qua preorder)} | \bigwedge \lambda(r) \leq \bigvee_{\pi(d)=r} \bigwedge \rho(d) \ (r \in R) \rangle$ .

In fact, we shall find it best to order not only  $G$ , but also  $R$  and  $D$ . The reasons for this will be revealed later. For the moment, let us say first, that we shall sometimes require the joins in the relations to be directed, and that requires an order on  $D$ . Second, in satisfying meet and join stability conditions, the generators  $G$  get incorporated into both  $R$  and  $D$  in such a way that the order cannot be neglected.

**5.7. DEFINITION.** *An ordered GRD-system is a system  $(G \text{ (qua poset)}, R, D, \lambda, \rho, \pi)$  in which  $G, R$  and  $D$  are equipped with partial orders  $\leq$ ,  $\pi$  is monotone and  $\lambda$  and  $\rho$  are monotone with respect to  $\leq_U$  on  $\mathcal{F}G$ .*

Notes:

- The orders on  $R$  and  $D$  make no difference to the frame that is presented.
- We expect the orders to be partial orders and not preorders.
- Ordinary GRD-systems can always be ordered by the discrete order.

**5.8. DEFINITION.**

1. Let  $(G, R, D, \lambda, \rho, \pi)$  be an ordered GRD-system in which  $G$  is a meet semilattice. Then by

$$(G \text{ (qua } \wedge \text{-semilattice)}, R, D, \lambda, \rho, \pi)$$

we denote the GRD-system

$$(G \text{ (qua poset)}, R', D', \lambda', \rho', \pi')$$

in which

$$\begin{aligned} R' &= R + 1 + G \times G \\ \lambda'(r) &= \lambda(r), \lambda'(*) = \emptyset, \lambda'(g_1, g_2) = \{g_1, g_2\} \\ D' &= D + 1 + G \times G \\ \rho'(d) &= \rho(d), \rho'(*) = \{\top_G\}, \rho'(g_1, g_2) = \{g_1 \wedge g_2\} \\ \pi' &= \pi + 1 + G \times G \end{aligned}$$

The extra relations are imposing  $\mathbf{true} \leq \top_G$  and  $(g_1) \wedge (g_2) \leq (g_1 \wedge g_2)$ .

2. Let  $(G, R, D, \lambda, \rho, \pi)$  be an ordered GRD-system in which  $G$  is a join semilattice. Then by

$$(G \text{ (qua } \vee\text{-semilattice)}, R, D, \lambda, \rho, \pi)$$

we denote the GRD-system

$$(G \text{ (qua poset)}, R', D', \lambda', \rho', \pi')$$

in which

$$\begin{aligned} R' &= R + 1 + G \times G \\ \lambda'(r) &= \lambda(r), \lambda'(*) = \{\perp_G\}, \lambda'(g_1, g_2) = \{g_1 \vee g_2\} \\ D' &= D + \emptyset + (G \times G + G \times G) \\ \rho'(d) &= \rho(d), \rho'((g_1, g_2)_1) = \{g_1\}, \rho'((g_1, g_2)_2) = \{g_2\} \\ \pi' &= \pi + ! + \nabla \end{aligned}$$

where  $\nabla$  is the codiagonal. The extra relations are imposing  $\perp_G \leq \mathbf{false}$  and  $(g_1 \vee g_2) \leq (g_1) \vee (g_2)$ .

## 6. Geometricity of internal directed joins

We have already mentioned in Section 2 that  $\mathbf{Loc}$  is  $\text{dcpo}$  enriched. However, the directed joins there are as indexed by *external* directed sets. More deeply we find that if  $(y_i)$  is a family of maps from  $X$  to  $Y$  indexed by a directed set  $I$  *internal to sheaves over*  $X$ , then they have a directed join  $\bigvee_{i \in I}^\uparrow y_i : X \rightarrow Y$ . Moreover, this is geometric: the directed joins are preserved under pullback along maps into  $X$ . (This corresponds at the global level to precomposition preserving directed joins, but we have to be more careful about looking after the directed index set.)

It is convenient to use the “ $\Omega$ -set” description of sheaves [FoScott 79]. This approach is slightly misleading, since the use of  $\Omega$ -sets is very specific to locales. It therefore obscures the fact that geometricity of directed joins for locales is just a special case of geometricity of filtered colimits for toposes. However, we retain the account here for its concreteness.

6.1. DEFINITION. *Let  $X$  be a locale. Then an  $\Omega X$ -set is a set  $S$  equipped with a function  $E : S^2 \rightarrow \Omega X$ ,  $(p, q) \mapsto E_{pq}$ , such that*

1.  $E_{pq} \leq E_{qp}$
2.  $E_{pq} \wedge E_{qr} \leq E_{pr}$

Working in  $\mathcal{S}X$ , we find that  $E$  defines a partial equivalence relation on  $S$  – condition (1) expresses symmetry, and condition (2) expresses transitivity. To put it another way, each point  $x$  of  $X$  gives a partial equivalence relation  $\equiv_x$  on  $S$ ,  $p \equiv_x q$  iff  $x \models E_{pq}$ , and in  $\mathcal{S}X$  we consider this for the generic point. Its quotient  $S/E$  is an object of  $\mathcal{S}X$ , i.e. a sheaf over  $X$ . (To form the quotient of a *partial* equivalence relation  $\equiv$  on  $S$ , you first restrict to  $S' = \{p \in S \mid p \equiv p\}$ , on which  $\equiv$  is an equivalence relation, and form  $S'/\equiv$ .) It is shown in [FoScott 79] that, with respect to suitable morphisms, the category of  $\Omega X$ -sets is equivalent to  $\mathcal{S}X$ , and so  $\Omega X$ -sets are just another way to describe sheaves over locales.

6.2. PROPOSITION. *Let  $X$  be a locale. Then a directed set at stage  $X$  is a set  $P$  equipped with a function  $L : P \times P \rightarrow \Omega X$ ,  $(p, q) \mapsto L_{pq}$ , that satisfies the following conditions.*

1.  $L_{pq} \leq L_{pp} \wedge L_{qq}$
2.  $L_{pq} \wedge L_{qr} \leq L_{pr}$
3.  $\mathbf{true} \leq \bigvee_{p \in P} L_{pp}$
4.  $L_{pp} \wedge L_{qq} \leq \bigvee_{r \in P} (L_{pr} \wedge L_{qr})$

PROOF. Let  $(P, E)$  be an  $\Omega X$ -set. As a binary relation, a partial order is given by a subsheaf of  $(P/E)^2$ , and this is given by a function  $L : P \times P \rightarrow \Omega X$ ,  $(p, q) \mapsto L_{pq}$ , satisfying

$$\begin{aligned} L_{pq} \wedge E_{p'p} \wedge E_{q'q} &\leq L_{p'q'} \\ L_{pq} &\leq E_{pp} \wedge E_{qq} \end{aligned}$$

The partial order properties (reflexive, transitive, antisymmetric) are then

$$\begin{aligned} E_{pp} &\leq L_{pp} \\ L_{pq} \wedge L_{qr} &\leq L_{pr} \\ L_{pq} \wedge L_{qp} &\leq E_{pq} \end{aligned}$$

and directedness is

$$\begin{aligned} \mathbf{true} &\leq \bigvee_{p \in P} E_{pp} \\ E_{pp} \wedge E_{qq} &\leq \bigvee_{r \in P} (L_{pr} \wedge L_{qr}) \end{aligned}$$

If we have  $E$  and  $L$  satisfying these, then we must have  $E_{pq} = L_{pq} \wedge L_{qp}$ . The  $\geq$  direction is already hypothesized, and for  $\leq$  we have

$$E_{pq} = E_{pp} \wedge E_{pp} \wedge E_{pq} \leq E_{pp} \wedge L_{pp} \wedge E_{pq} \leq L_{pq}$$

and  $E_{pq} \leq L_{qp}$  by symmetry. It follows that  $E_{pp} = L_{pp}$ , and we can easily deduce the properties listed in the statement.

Conversely, suppose we have  $L$  satisfying the stated properties and we define  $E_{pq} = L_{pq} \wedge L_{qp}$ . We can then check that  $(P, E)$  is an  $\Omega X$ -set and that  $L$  defines a directed partial order on it. ■

**6.3. PROPOSITION.** *Let  $X$  and  $Y$  be locales and let  $P$  be a directed set at stage  $X$  (equipped with order  $L$ ). Then a  $P$ -indexed family of points of  $Y$  is a family of maps  $y_p : L_{pp} \rightarrow Y$  such that  $y_p|_{L_{pq}} \sqsubseteq y_q|_{L_{pq}} : L_{pq} \rightarrow Y$ .*

**PROOF.** Each element of  $P/L$  is represented by an element  $p$  of  $P$ , but it is only partial – it exists where  $L_{pp}$  holds. The corresponding point of  $Y$  is a map  $y_p : L_{pp} \rightarrow Y$ . The remaining condition represents the monotonicity of the family, if  $p \leq q$  then  $y_p \sqsubseteq y_q$ . ■

**6.4. THEOREM.** *Let  $X$  and  $Y$  be locales, let  $P$  be a directed set over  $X$ , and let  $(y_p)$  be a  $P$ -indexed family of points of  $Y$  at stage  $X$ . Then the directed join  $\bigsqcup_{p \in P}^\uparrow y_p$  is the map  $y : X \rightarrow Y$  defined by*

$$\Omega y(b) = \bigvee_{p \in P} \Omega y_p(b)$$

(Note that we are taking  $\Omega L_{pp} = \downarrow L_{pp} \subseteq \Omega X$ .)

**PROOF.** The requirement on  $y$  is that it should be the least map from  $X$  to  $Y$  such that for every  $p \in P$  we have

$$y_p \sqsubseteq y|_{L_{pp}} : L_{pp} \rightarrow Y,$$

in other words that  $\Omega y_p(b) \leq \Omega y(b)$  for every  $p \in P$  and every  $b \in \Omega Y$ . The function  $\Omega y$  in the statement is clearly the least *function* from  $\Omega Y$  to  $\Omega X$  that satisfies this, and so it suffices to show that it is a frame homomorphism.

It is evident that  $\Omega y$  preserves all joins. For finite meets, first,

$$\Omega y(\mathbf{true}) = \bigvee_{p \in P} \Omega y_p(\mathbf{true}) = \bigvee_{p \in P} L_{pp} = \mathbf{true}$$



For binary meets,

$$\begin{aligned}
 \Omega y(b) \wedge \Omega y(c) &= \bigvee_{p \in P} \Omega y_p(b) \wedge \bigvee_{q \in P} \Omega y_q(c) \\
 &= \bigvee_{p, q \in P} (\Omega y_p(b) \wedge \Omega y_q(c)) \\
 &= \bigvee_{p, q \in P} (\Omega y_p(b) \wedge \Omega y_q(c) \wedge L_{pp} \wedge L_{qq}) \\
 &\leq \bigvee_{p, q, r \in P} (\Omega y_p(b) \wedge \Omega y_q(c) \wedge L_{pr} \wedge L_{qr}) \\
 &\leq \bigvee_{p, q, r \in P} (\Omega y_r(b) \wedge \Omega y_r(c) \wedge L_{pr} \wedge L_{qr}) \\
 &\leq \bigvee_{r \in P} \Omega y_r(b \wedge c) = \Omega y(b \wedge c)
 \end{aligned}$$

■

Next we prove that composition of maps distributes on both sides over directed joins. For precomposition, we must explain how to shift the stage of a directed set.

6.5. PROPOSITION. *Let  $X$  be a locale, let  $P$  be a directed set at stage  $X$ , and let  $f : W \rightarrow X$  be a map. Then the pullback of  $P$  along  $f$ ,  $f^*P$ , is the directed set at stage  $W$  with the same underlying set  $P$  and with order  $f^*L$  defined by*

$$(f^*L)_{pq} = \Omega f(L_{pq}) = f^*(L_{pq})$$

PROOF. This is simply the way  $f^*$  works with  $\Omega$ -sets. ■

6.6. DEFINITION. *Let  $X$  be a locale, let  $P$  be a directed set at stage  $X$ , and let  $f : W \rightarrow X$  be a map. If  $(y_p)$  is a  $P$ -indexed family of points of  $Y$ , then we denote by  $(f; y_p)$  the following  $f^*P$ -indexed family of points of  $Y$ . For each  $p$ , let  $i_p : L_{pp} \hookrightarrow X$  be the sublocale inclusion. Then, abusing notation slightly,  $f; y_p$  here denotes  $(i_p^*f); y_p : f^*L_{pp} \rightarrow L_{pp} \rightarrow Y$ .*

6.7. THEOREM. *Let  $X$  and  $Y$  be locales, let  $P$  be a directed set over  $X$ , and let  $(y_p)$  be a  $P$ -indexed family of points of  $Y$  at stage  $X$ .*

1. *If  $f : W \rightarrow X$  is a map, then*

$$f; \bigsqcup_{p \in P}^\uparrow y_p = \bigsqcup_{p \in f^*P}^\uparrow (f; y_p)$$

2. *If  $g : Y \rightarrow Z$  is a map, then*

$$(\bigsqcup_{p \in P}^\uparrow y_p); g = \bigsqcup_{p \in P}^\uparrow (y_p; g)$$

PROOF. 1. We have

$$\begin{aligned} \Omega(f; \bigsqcup_{p \in P}^\uparrow y_p)(b) &= \Omega f(\bigvee_p \Omega y_p(b)) = \bigvee_p \Omega f(\Omega y_p(b)) \\ \Omega(\bigsqcup_{p \in f^*P}^\uparrow (f; y_p))(b) &= \bigvee_p \Omega(f; y_p)(b) \end{aligned}$$

We must therefore show that  $\Omega f(\Omega y_p(b)) = \Omega(f; y_p)(b)$  – this is not completely obvious, because we have slightly redefined what  $f; y_p$  means. If  $i_p : L_{pp} \hookrightarrow X$  is the sublocale inclusion then  $\Omega(f; y_p)(b) = \Omega(i_p^* f)(\Omega y_p(b))$ , where  $\Omega y_p(b) \in \Omega L_{pp}$ . But when we identify this with  $\downarrow L_{pp}$  in  $\Omega X$ , and make a similar identification for  $f^* L_{pp}$ , we find this is equal to  $\Omega f(\Omega y_p(b))$ .

2.

$$\begin{aligned} \Omega((\bigsqcup_{p \in P}^\uparrow y_p); g)(c) &= \bigvee_p \Omega y_p(\Omega g(c)) \\ &= \bigvee_p \Omega(y_p; \Omega g)(c) = \Omega(\bigsqcup_{p \in P}^\uparrow (y_p; g))(c) \end{aligned}$$

■

## 7. Geometricity of the lower powerlocale

We shall now use the technique of Proposition 5.5 to prove that the powerlocales  $P_U$  and  $P_L$ , and hence also the double powerlocale  $\mathbb{P}$ , are geometric. The argument makes essential use of the *coverage theorems*, by which frame presentations are transformed into suplattice or preframe presentations.

We deal first with the suplattice coverage theorem, which is Johnstone’s original form [Johnstone 82].

**7.1. THEOREM.** [Suplattice coverage theorem] *Let  $S$  be a meet semilattice, and let  $C \subseteq S \times \mathcal{P}S$  have the following properties. (If  $(x, U) \in C$  then we say that  $U$  covers  $x$ .)*

- *If  $U$  covers  $x$  and  $u \in U$  then  $u \leq x$ .*
- *(Meet stability) If  $U$  covers  $x$  and  $y \in S$  then  $\{y \wedge u \mid u \in U\}$  covers  $y \wedge x$ .*

*Then*

$$\begin{aligned} \text{Fr}\langle S \text{ (qua meet semilattice)} \mid x \leq \bigvee U \text{ (} U \text{ covers } x) \rangle \\ \cong \text{SupL}\langle S \text{ (qua poset)} \mid x \leq \bigvee U \text{ (} U \text{ covers } x) \rangle \end{aligned}$$

PROOF. [Johnstone 82] proved that the frame presented on the left can be represented concretely as the lattice of “ $C$ -ideals”. It was proved in [AbrVic 93] that this lattice could be given the suplattice presentation on the right. The arguments used there are constructive. ■

7.2. DEFINITION. A site is an ordered GRD-system of the form  $(G \text{ (qua } \wedge\text{-semilattice)}, R, D)$  equipped with the following structure and properties.

1. Factorizations of  $\lambda$  and  $\rho$  via the singleton embedding  $\{-\} : G \rightarrow \mathcal{F}G$ ,  $\lambda(r) = \{\lambda_0(r)\}$  and  $\rho(d) = \{\rho_0(d)\}$ .
2.  $\rho_0(d) \leq \lambda_0(\pi(d))$ .
3. Actions of  $G$  on  $R$  and  $D$ , denoted  $(g, r) \mapsto g \cdot r$  and  $(g, d) \mapsto g \cdot d$ .
4. The functions  $\cdot : G \times R \rightarrow R$  and  $\cdot : G \times D \rightarrow D$  are both monotone.
5.  $\pi(g \cdot d) = g \cdot \pi(d)$
6. If  $\pi(d') = g \cdot r$  then there is some  $d$  such that  $\pi(d) = r$  and  $d' \leq g \cdot d$ .
7.  $\lambda_0(g \cdot r) = g \wedge \lambda_0(r)$
8.  $\rho_0(g \cdot r) = g \wedge \rho_0(r)$

This is mostly just a formalization in terms of GRD-systems of the conditions used in Theorem 7.1, so it is not surprising that we now have the following.

7.3. COROLLARY. Let  $(G, R, D)$  be a site, and let  $X$  be presented by  $(G \text{ (qua } \wedge\text{-semilattice)}, R, D)$ . Then  $P_L X$  is presented by  $(G \text{ (qua poset)}, R, D)$ .

PROOF. We define a coverage relation  $C \subseteq G \times \mathcal{P}G$  whose elements are the pairs

$$(\lambda_0(r), \{\rho_0(d) \mid \pi(d) = r\})$$

for  $r \in R$ . The fact that it is a coverage comes from the structure of the site; in particular the  $G$ -action on  $R$  and  $D$  gives meet stability. Then  $\Omega X$  is

$$\text{Fr}\langle G \text{ (qua } \wedge\text{-semilattice)} \mid x \leq \bigvee U \text{ (} U \text{ covers } x) \rangle$$

Using Theorem 7.1 we see that

$$\begin{aligned} \Omega P_L X &\cong \text{Fr}\langle \Omega X \text{ (qua suplattice)} \rangle \\ &\cong \text{Fr}\langle G \text{ (qua poset)} \mid x \leq \bigvee U \text{ (} U \text{ covers } x) \rangle \end{aligned}$$

and this is presented by  $(G \text{ (qua poset)}, R, D)$ . ■

In the light of this, geometricity of the lower powerlocale reduces to showing that any presentation can be geometrically transformed into an equivalent site. First, however, we take the opportunity to prove a Lemma, probably well-known, that will be useful in various places.

#### 7.4. LEMMA.

1. Let  $P$  be a poset. Then  $\mathcal{F}P/\leq_L$  is a free join semilattice over  $P$  (qua poset).
2. Let  $P$  be a poset. Then  $\mathcal{F}P/\leq_U$  is a free meet semilattice over  $P$  (qua poset).
3. Let  $P$  be a meet semilattice. Then  $\mathcal{F}P/\leq_L$  is a free distributive lattice over  $P$  (qua meet semilattice).

PROOF. 1.  $\emptyset$  and  $\cup$  in  $\mathcal{F}P$  provide (nullary and binary) joins with respect to  $\leq_L$ . For instance,  $U \leq_L U \cup V$ , and if  $U \leq_L W$  and  $V \leq_L W$  then  $U \cup V \leq_L W$ . The singleton embedding of  $P$  in  $\mathcal{F}P$ ,  $x \mapsto \{x\}$ , is monotone with respect to  $\leq_L$ : if  $x \leq y$  then  $\{x\} \leq_L \{y\}$ .

Now let  $S$  be a join semilattice and  $f : P \rightarrow S$  monotone. This extends uniquely to a semilattice homomorphism  $f_1 : \mathcal{F}P \rightarrow S$ ,  $U \mapsto \bigvee \{f(u) \mid u \in U\}$ , and this factors via  $\mathcal{F}P/\leq_L$ : if  $U \leq_L V$  then  $f_1(U) \leq f_1(V)$ .

2. This follows by duality from the fact that  $S \leq_U T$  iff  $T(\geq)_L S$ . Hence a meet semilattice homomorphism from  $\mathcal{F}P/\leq_U$  to  $S$  is equivalent to a join semilattice homomorphism from  $\mathcal{F}(P^{\text{op}})/\leq_L$  to  $S^{\text{op}}$ , which is equivalent to a monotone function from  $P^{\text{op}}$  to  $S^{\text{op}}$  and hence to a monotone function from  $P$  to  $S$ .

3. The top element is  $\{\top_P\}$ , and binary meets are defined by

$$U \wedge V = \{u \wedge v \mid u \in U, v \in V\}$$

We have  $U \wedge V \leq_L U$  and if  $W \leq_L U$  and  $W \leq_L V$  then  $W \leq_L U \wedge V$ . For distributivity we have

$$U \wedge (V \cup W) \leq_L (U \wedge V) \cup (U \wedge W)$$

Suppose  $S$  is a distributive lattice and  $f : P \rightarrow S$  a meet semilattice homomorphism. Then the unique extension to a join semilattice homomorphism  $f_2 : \mathcal{F}P/\leq_L \rightarrow S$  is in fact a lattice homomorphism, for

$$\begin{aligned} f_1(U \wedge V) &= f_1(\{u \wedge v \mid u \in U, v \in V\}) \\ &= \bigvee \{f(u) \wedge f(v) \mid u \in U, v \in V\} \\ &= f_1(U) \wedge f_1(V) \end{aligned}$$

■

If  $U \in \mathcal{FP}$  then we shall write  $[U]$  for its image in  $\mathcal{FP}/\leq_U$  or  $\mathcal{FP}/\leq_L$  as appropriate.

7.5. THEOREM. *The lower powerlocale  $P_L$  is geometric.*

1. *If  $X$  is presented by an ordered GRD-system  $(G$  (qua poset),  $R, D$ ), then  $P_L X$  is presented by a system  $(G'$  (qua poset),  $R', D')$  where*

$$\begin{aligned} G' &= \mathcal{F}G/\leq_U \\ R' &= G' \times R \\ \lambda(S, r) &= \{S \wedge \lambda(r)\} \\ D' &= G' \times D \\ \rho(S, d) &= \{S \wedge \rho(d) \wedge \lambda(\pi(d))\} \\ \pi &= G' \times \pi \end{aligned}$$

2.  $\Omega \downarrow$  takes  $[U] \mapsto \bigwedge U$ ,  $\Omega \sqcup$  takes  $[U] \mapsto [\{[U]\}]$ .
3. *Let  $X_i$  be presented by  $(G_i$  (qua poset),  $R_i, D_i)$  ( $i = 1, 2$ ). Let  $f : X_1 \rightarrow X_2$  be a map, and suppose  $\phi \subseteq \mathcal{F}G_1 \times G_2$  is such that, for  $g \in G_2$ ,*

$$\Omega f(g) = \bigvee \{ \bigwedge S \mid (S, g) \in \phi \}.$$

Let

$$\phi' = \{ (\{ \bigcup_{(S,g) \in \phi_0} S \}, V) \mid \phi_0 \in \mathcal{F}\phi, \mathcal{F} \text{snd } \phi_0 = V \} \subseteq \mathcal{F}\mathcal{F}G_1 \times \mathcal{F}G_2.$$

Then  $P_L f$  is similarly described by  $\phi'$ ,

$$\Omega P_L f([V]) = \bigvee_{U \in \mathcal{U}} \{ \bigwedge [U] \mid (U, V) \in \phi' \}.$$

PROOF. 1. If  $U \in \mathcal{F}G$ , let us write  $[U]$  for its image in  $G'$ .

We have

$$\begin{aligned} \Omega X &= \text{Fr} \langle G \text{ (qua poset)} \mid \bigwedge \lambda(r) \leq \bigvee_{\pi(d)=r} \bigwedge \rho(d) \ (r \in R) \rangle \\ &\cong \text{Fr} \langle G' \text{ (qua } \wedge \text{-semilattice)} \mid \\ &\quad S \wedge [\lambda(r)] \leq \bigvee_{\pi(d)=r} (S \wedge [\rho(d)] \wedge [\lambda(r)]) \ (r \in R, S \in G') \rangle \end{aligned}$$

This is shown by demonstrating the mutually inverse frame homomorphisms.  $\alpha$  from left to right is defined by  $\alpha(g) = [\{g\}]$ ,  $\beta$  from right to left by  $\beta([U]) = \bigwedge U$ , using Lemma 7.4 (2).

We must still check that the relations are respected, but once that is done the two homomorphisms will be mutually inverse. On one side,  $\beta \circ \alpha(g) = \beta(\{\{g\}\}) = g$ . On the other we have

$$\alpha \circ \beta(\{\{g\}\}) = \alpha(g) = \{\{g\}\}$$

after which  $\alpha \circ \beta = \text{Id}$  by Lemma 7.4.

Now we check that the relations are respected. From left to right, there is only one. We have

$$\begin{aligned} \alpha(\bigwedge \lambda(r)) &= \bigwedge_{g \in \lambda(r)} \{\{g\}\} = [\bigcup_{g \in \lambda(r)} \{g\}] = [\lambda(r)] \\ &\leq \bigvee_{\pi(d)=r} [\rho(d)] = \bigvee_{\pi(d)=r} \bigwedge_{g \in \rho(d)} \{\{g\}\} \\ &= \alpha(\bigvee_{\pi(d)=r} \bigwedge \rho(d)). \end{aligned}$$

Going from right to left, we have

$$\begin{aligned} \beta(S \wedge [\lambda(r)]) &= \beta(S) \wedge \bigwedge \lambda(r) \\ &\leq \beta(S) \wedge \bigvee_{\pi(d)=r} \bigwedge \rho(d) \wedge \bigwedge \lambda(r) \\ &= \bigvee_{\pi(d)=r} (\beta(S) \wedge \bigwedge \rho(d) \wedge \bigwedge \lambda(r)) \\ &= \beta(\bigvee_{\pi(d)=r} (S \wedge [\rho(d)] \wedge [\lambda(r)])) \end{aligned}$$

The right-hand presentation corresponds to the GRD-system described, and moreover it is a site, with  $G'$ -actions given by

$$\begin{aligned} S \cdot (T, r) &= (S \wedge T, r) \\ S \cdot (T, d) &= (S \wedge T, d) \end{aligned}$$

The result now follows from Corollary 7.3.

2.  $U \in \mathcal{F}G$  represents  $\diamond \bigwedge_{g \in U} g \in \Omega \text{P}_L X$ . Under  $\Omega \downarrow$ , this maps to  $\bigwedge U$ . Under  $\Omega \sqcup$ , it maps to  $\diamond \diamond \bigwedge_{g \in U} g$ , which corresponds to  $\{\{U\}\}$ .

3. Under  $\Omega \text{P}_L f$ ,  $\diamond \bigwedge_{g \in V} g$  maps to

$$\diamond \bigwedge_{g \in V} \bigvee_{(S,g) \in \phi} \bigwedge S$$

If  $V \in \mathcal{F}G$ , let us write

$$\begin{aligned} \Phi(V) &= \{\phi_0 \in \mathcal{F}\phi \mid \forall (S, g) \in \phi_0. g \in V \text{ and } \forall g \in V. \exists (S, g) \in \phi_0\} \\ &= \{\phi_0 \in \mathcal{F}\phi \mid \mathcal{F} \text{snd}(\phi_0) = V\} \end{aligned}$$

We must therefore show that

$$\bigwedge_{g \in V} \bigvee_{(S,g) \in \phi} \bigwedge S = \bigvee_{\phi_0 \in \Phi(V)} \bigwedge_{(S,g) \in \phi_0} \bigwedge S$$

For  $\geq$ , suppose we have  $\phi_0 \in \Phi(V)$ , and we have  $h \in V$ . We can find  $(T, h) \in \phi_0$  for some  $T$ , so

$$\bigwedge_{(S,g) \in \phi_0} \bigwedge S \leq \bigwedge T \leq \bigvee_{(S,h) \in \phi} \bigwedge S$$

For  $\leq$ , we use induction on  $V$  (see [Vickers 99]). If  $V = \emptyset$ , then the LHS is **true**. For the RHS, the only possible  $\phi_0$  is  $\emptyset$  and it follows that the RHS also is **true**. For the induction step, suppose  $V = V' \cup \{h\}$ .

$$\begin{aligned} \text{LHS} &= \bigwedge_{g \in V'} \bigvee_{(S,g) \in \phi} \bigwedge S \wedge \bigvee_{(T,h) \in \phi} \bigwedge T \\ &= \bigvee_{\phi_0 \in \Phi(V')} \bigwedge_{(S,g) \in \phi_0} \bigwedge S \wedge \bigvee_{(T,h) \in \phi} \bigwedge T \\ &= \bigvee \{ \bigwedge_{(S,g) \in \phi_0 \cup \{(T,h)\}} \bigwedge S \mid \phi_0 \in \Phi(V'), (T,h) \in \phi \} \\ &\leq \bigvee_{\phi_0 \in \Phi(V)} \bigwedge_{(S,g) \in \phi_0} \bigwedge S \end{aligned}$$

Hence

$$\diamond \bigwedge_{g \in V} \bigvee_{(S,g) \in \phi} \bigwedge S = \bigvee_{\phi_0 \in \Phi(V)} \diamond \bigwedge_{(S,g) \in \phi_0} \bigwedge S = \bigvee_{\phi_0 \in \Phi(V)} \diamond \bigwedge ( \bigcup_{(S,g) \in \phi_0} S )$$

which gives us the result. ■

## 8. Geometricity of the upper powerlocale

Turning now to the upper powerlocale, we first define a preframe analogue of site.

**8.1. DEFINITION.** *An ordered GRD-system  $(G$  (qua  $\vee$ -semilattice),  $R, D$ ) is a preframe site if it has the following structure and properties.*

1. For each  $r$  in  $R$ , the set  $D_r = \{d \in D \mid \pi(d) = r\}$  is directed under  $\leq$ .
2. For each  $d$  in  $D$ ,  $\rho(d) \leq_U \lambda(\pi(d))$ .
3. Actions of  $G$  on  $R$  and  $D$ , denoted  $(a, r) \mapsto a * r$  and  $(a, d) \mapsto a * d$ .
4. The functions  $*$  :  $G \times R \rightarrow R$  and  $*$  :  $G \times D \rightarrow D$  are both monotone.
5.  $\pi(a * d) = a * \pi(d)$ .

- 6. If  $\pi(d') = g * r$  then there is some  $d$  such that  $\pi(d) = r$  and  $d' \leq g * d$ .
- 7.  $\lambda(a * r) \equiv_U \{a \vee b \mid b \in \lambda(r)\}$ .
- 8.  $\rho(a * d) \equiv_U \{a \vee b \mid b \in \rho(d)\}$ .

Conditions 1 and 2 here say that the relations are all of the form

$$\bigwedge S \leq \bigvee_i^\uparrow \bigwedge T_i$$

with a directed join on the right, such that each  $\bigwedge T_i \leq \bigwedge S$ . Conditions 3-8 express a *join stability* property, namely that if we have such a relation in  $R$ , and  $a \in G$ , then we also have the relation

$$\bigwedge \{a \vee b \mid b \in S\} \leq \bigvee_i^\uparrow \bigwedge \{a \vee b \mid b \in T_i\}$$

8.2. THEOREM. [Preframe coverage theorem] *Let  $(G, R, D)$  be a preframe site. Then*

$$\begin{aligned} \text{Fr}\langle G \text{ (qua join semilattice)} \mid \bigwedge \lambda(r) \leq \bigvee_{\pi(d)=r}^\uparrow \bigwedge \rho(d) \text{ (} r \in R \text{)} \rangle \\ \cong \text{PreFr}\langle G \text{ (qua poset)} \mid \bigwedge \lambda(r) \leq \bigvee_{\pi(d)=r}^\uparrow \bigwedge \rho(d) \text{ (} r \in R \text{)} \rangle \end{aligned}$$

PROOF. [JoVic 91] The first task is to show that the preframe presented on the right actually exists. Once that is done, binary joins can be defined as preframe bilinear maps, and proved to be joins. ■

8.3. COROLLARY. *Let  $(G, R, D)$  be a preframe site, and let  $X$  be presented by  $(G \text{ (qua } \vee\text{-semilattice)}, R, D)$ . Then  $P_U X$  is presented by  $(G \text{ (qua poset)}, R, D)$ .*

We shall now show that the upper powerlocale  $P_U$  is geometric. Just as for  $P_L$ , the strategy is to show how to manipulate an arbitrary presentation into an equivalent preframe site. However, in doing so we shall also need to use finite distributivity and we must first say something about the geometric nature of that.

Let  $X$  be a distributive lattice and  $\mathcal{V}$  a finite set of finite subsets of  $X$ . As ever, “finite” here means Kuratowski finite, so  $\mathcal{V} \in \mathcal{FF}X$ . We are interested in the problem of writing the join of meets

$$\bigvee_{V \in \mathcal{V}} \bigwedge V$$

as a meet of joins. If one imagines writing this out, one finds that each of the joins has a disjunct from each  $V$  in  $\mathcal{V}$ , so

$$\bigvee_{V \in \mathcal{V}} \bigwedge V = \bigwedge \{ \bigvee_{V \in \mathcal{V}} \gamma(V) \mid \gamma \text{ a choice function for } \mathcal{V} \}$$

A choice function for  $\mathcal{V}$  is a function  $\gamma : \mathcal{V} \rightarrow \bigcup \mathcal{V}$  such that  $\gamma(V) \in V$  for every  $V$ . Since  $\mathcal{V}$  is finite, the graph of  $\gamma$  must be a finite subset of  $\mathcal{V} \times \bigcup \mathcal{V}$ , but unfortunately



the set of such  $\gamma$ 's is not a geometrically definable subset of  $\mathcal{F}(\mathcal{V} \times \bigcup \mathcal{V})$ . This is because the single-valuedness property of a function,  $\forall V, v, v'. (V\gamma v \wedge V\gamma v' \rightarrow v = v')$ , is not geometric. (To make it so we should need decidable equality on  $X$ , when the property can be expressed as  $\forall (V, v) \in \gamma. \forall (V', v') \in \gamma. (V \neq V' \vee v = v')$ .) However, it turns out that single-valuedness is not needed for our applications. We can make do with choice relations that are total but not necessarily single-valued.

8.4. DEFINITION. *Let  $X$  be a set and  $\mathcal{V} \in \mathcal{F}\mathcal{F}X$ . A choice for  $\mathcal{V}$  is some  $\gamma \in \mathcal{F}(\mathcal{V} \times \bigcup \mathcal{V})$  satisfying*

- $\forall V \in \mathcal{V}. \exists v \in V. (V, v) \in \gamma$
- $\forall (V, v) \in \gamma. v \in V$

8.5. DEFINITION. *We write  $\text{Ch}(\mathcal{V})$  for the set of choices of  $\mathcal{V}$ , and  $\text{Im } \gamma$  for the image of  $\gamma$ , i.e. its direct image under the second projection to  $\bigcup \mathcal{V}$  (note that  $\text{Im } \gamma$  is finite).*

8.6. PROPOSITION. *Let  $X$  be a set and  $\mathcal{V} \in \mathcal{F}\mathcal{F}X$ . Then  $\text{Ch}(\mathcal{V})$  is finite.*

PROOF. We use the “ $\mathcal{F}$ -recursion” of [Vickers 99] to define a function  $\text{Ch} : \mathcal{F}\mathcal{F}X \rightarrow \mathcal{F}\mathcal{F}(\mathcal{F}X \times X)$  and show that  $\gamma \in \mathcal{F}(\mathcal{F}X \times X)$  is in  $\text{Ch}(\mathcal{V})$  iff it is a choice of  $\mathcal{V}$  as defined above.

First,  $\text{Ch}(\emptyset) = \{\emptyset\}$ .  $\emptyset$  is, vacuously, a choice of  $\emptyset$ , and it is the only finite subset of  $\emptyset \times \bigcup \emptyset$ .

Next,

$$\text{Ch}(\mathcal{V} \cup \{U\}) = \{\gamma \cup (\{U\} \times S) \mid \gamma \in \text{Ch}(\mathcal{V}) \text{ and } S \in \mathcal{F}^+U\}$$

( $\mathcal{F}^+U$  denotes the set of nonempty finite subsets of  $U$  – note that for finite sets, emptiness is decidable.) There are two proof obligations to be discharged in showing that this definition is good. For  $\mathcal{W} \in \mathcal{F}\mathcal{F}(\mathcal{F}X \times X)$  and  $U \in \mathcal{F}X$  let us write

$$\Phi(\mathcal{W}, U) = \{\gamma \cup (\{U\} \times S) \mid \gamma \in \mathcal{W} \text{ and } S \in \mathcal{F}^+U\}$$

so we are attempting to define  $\text{Ch}$  to satisfy

$$\text{Ch}(\mathcal{V} \cup \{U\}) = \Phi(\text{Ch } \mathcal{V}, U)$$

Then we must show

$$\begin{aligned} \Phi(\Phi(\mathcal{W}, U), V) &= \Phi(\Phi(\mathcal{W}, V), U) \\ \Phi(\Phi(\mathcal{W}, U), U) &= \Phi(\mathcal{W}, U) \end{aligned}$$

Both of these are easy and we deduce by  $\mathcal{F}$ -recursion that  $\text{Ch}$  is well-defined. Now for correctness, we can suppose by induction that  $\text{Ch}(\mathcal{V})$  is correct (i.e. the set of choices of  $\mathcal{V}$ ). Then clearly all the elements of  $\text{Ch}(\mathcal{V} \cup \{U\})$  are choices of  $\mathcal{V} \cup \{U\}$ . Conversely,

let  $\gamma$  be a choice of  $\mathcal{V} \cup \{U\}$ , and write  $\gamma = \gamma_1 \cup \gamma_2$  where the  $\gamma_i$ s are both finite, for every  $(V, v)$  in  $\gamma_1$  we have  $V$  in  $\mathcal{V}$ , and for every  $(V, v)$  in  $\gamma_2$  we have  $V = U$ . (If  $U \in \mathcal{V}$  then  $\gamma_1$  and  $\gamma_2$  are not uniquely defined. Because this possibility cannot be eliminated geometrically, the reasoning is slightly more elaborate than one might have expected.) We have  $\forall V \in \mathcal{V}. \exists v \in V. (V, v) \in \gamma$ , and so there is a finite total relation  $R$  from  $\mathcal{V}$  to  $\bigcup \mathcal{V}$  contained in  $\gamma$ . Replacing  $\gamma_1$  by  $\gamma_1 \cup R$  we can suppose that  $\gamma_1$  is total on  $\mathcal{V}$ . We also have some  $u \in U$  such that  $(U, u) \in \gamma$ , so replacing  $\gamma_2$  by  $\gamma_2 \cup \{(U, u)\}$  we can suppose that  $\gamma_2$  is of the form  $\{U\} \times S$  for some finite non-empty set  $S$ . Hence  $\gamma \in \text{Ch}(\mathcal{V} \cup \{U\})$ . ■

8.7. THEOREM. *Let  $L$  be a distributive lattice and let  $\mathcal{V} \in \mathcal{FFL}$ . Then*

$$\bigvee_{V \in \mathcal{V}} \bigwedge V = \bigwedge_{\gamma \in \text{Ch}(\mathcal{V})} \bigvee \text{Im } \gamma.$$

PROOF. By  $\mathcal{F}$ -induction on  $\mathcal{V}$ . (See [Vickers 99]. Note that this is not induction on the cardinality of  $\mathcal{V}$ , which in general does not exist. Rather, it is induction on the length of an enumeration of  $\mathcal{V}$ .) If  $\mathcal{V}$  is empty, then the LHS is  $\perp_L$  and the RHS is  $\bigwedge_{\gamma \in \{\emptyset\}} \bigvee \emptyset = \perp_L$ . For the induction step,

$$\begin{aligned} \bigvee_{V' \in \{V\} \cup \mathcal{V}} \bigwedge V' &= \bigwedge V \vee \bigvee_{V' \in \mathcal{V}} \bigwedge V' = \bigwedge V \vee \bigwedge_{\gamma' \in \text{Ch}(\mathcal{V})} \bigvee \text{Im } \gamma' \\ &= \bigwedge_{v \in V, \gamma' \in \text{Ch}(\mathcal{V})} (v \vee \bigvee \text{Im } \gamma') = \bigwedge_{\gamma \in \text{Ch}(\mathcal{V})} \bigvee \text{Im } \gamma \end{aligned}$$

■

Of course, there is a dual result in which meets and joins are interchanged.

8.8. THEOREM. *The upper powerlocale  $P_U$  is geometric.*

1. *If  $X$  is presented by an ordered GRD-system  $(G$  (qua poset),  $R, D)$ , then  $P_U X$  is presented by a system  $(G'$  (qua poset),  $R', D')$  where*

$$\begin{aligned} G' &= \mathcal{F}G / \leq_L \\ R' &= G' \times R \\ \lambda(S, r) &= \{S \vee [\{g\}] \mid g \in \lambda(r)\} \\ D' &= \{(S, D_0, r) \in G' \times \mathcal{F}D \times R \mid \forall d \in D_0. \pi(d) = r\} / (\leq \times \leq_L \times \leq) \\ \rho(S, D_0, r) &= \{(S \vee [\text{Im } \gamma]) \mid \gamma \in \text{Ch}\{\lambda(r) \cup \rho(d) \mid d \in D_0\}\} \\ \pi(S, D_0, r) &= (S, r) \end{aligned}$$

2.  $\Omega \uparrow$  takes  $[U] \mapsto \bigvee U$ ,  $\Omega \sqcap$  takes  $[U] \mapsto [\{[U]\}]$ .

3. Let  $X_i$  be presented by  $(G_i \text{ (qua poset)}, R_i, D_i)$  ( $i = 1, 2$ ). Let  $f : X_1 \rightarrow X_2$  be a map, and suppose  $\phi \subseteq \mathcal{F}G_1 \times G_2$  is such that, for  $g \in G_2$ ,

$$\Omega f(g) = \bigvee \{ \bigwedge S \mid (S, g) \in \phi \}.$$

Let

$$\begin{aligned} \phi' &= \{ (\{ \text{Im } \gamma \mid \gamma \in \text{Ch}(\{ S \mid (S, g) \in \phi_0 \}) \}, V) \mid \phi_0 \in \mathcal{F}\phi, \mathcal{F} \text{snd } \phi_0 \subseteq V \} \\ &\subseteq \mathcal{F}\mathcal{F}G_1 \times \mathcal{F}G_2 \end{aligned}$$

Then  $P_U f$  is similarly described by  $\phi'$ ,

$$P_U f([V]) = \bigvee_{U \in \mathcal{U}} \{ \bigwedge [U] \mid (U, V) \in \phi' \}$$

PROOF. 1. Note that  $(G' \text{ (qua } \vee \text{-semilattice)}, R', D')$  is a preframe site. To show  $\rho$  is well-defined, suppose  $S \leq S', D_0 \leq_L D'_0$  and  $r \leq r'$ . We must show that  $\rho(S, D_0, r) \leq_U \rho(S', D'_0, r')$ . Suppose

$$\gamma' \in \text{Ch} \{ \lambda(r') \cup \rho(d') \mid d' \in D'_0 \}.$$

For every  $d \in D_0$  we can find  $d' \in D'_0$  with  $d \leq d'$ . Then

$$\lambda(r) \cup \rho(d) \leq_U \lambda(r') \cup \rho(d').$$

We can find  $g' \in \lambda(r') \cup \rho(d')$  with  $(d', g') \in \gamma'$  and then  $g \in \lambda(r) \cup \rho(d)$  with  $g \leq g'$ . It follows that there is some  $\gamma \in \text{Ch} \{ \lambda(r) \cup \rho(d) \mid d \in D_0 \}$  such that if  $(d, g) \in \gamma$ , then there is some  $(d', g') \in \gamma'$  such that  $d \leq d'$  and  $g \leq g'$ . Then  $S \vee \text{Im } \gamma \leq S' \vee \text{Im } \gamma'$ , because  $\text{Im } \gamma \leq_L \text{Im } \gamma'$ .

To show condition 2 in Definition 8.1 we must show  $\rho(S, D_0, r) \leq_U \lambda(S, r)$ . If  $g \in \lambda(r)$ , so  $S \vee [\{g\}] \in \lambda(S, r)$ , then  $D_0 \times \{g\}$  is a constant choice for  $\{ \lambda(r) \cup \rho(d) \mid d \in D_0 \}$  and its image is  $\{g\}$ . Hence  $S \vee [\{g\}] \in \rho(S, D_0, r)$ .

The  $G'$ -actions are given by

$$\begin{aligned} S * (T, r) &= (S \vee T, r) \\ S * (T, D_0, r) &= (S \vee T, D_0, r) \end{aligned}$$

Once we have shown that it is equivalent to  $(G \text{ (qua poset)}, R, D)$ , we can apply Corollary 8.3.

We have

$$\begin{aligned} \Omega X &= \text{Fr} \langle G \text{ (qua poset)} \mid \bigwedge \lambda(r) \leq \bigvee_{\pi(d)=r} \bigwedge \rho(d) \ (r \in R) \rangle \\ &\cong \text{Fr} \langle G \text{ (qua poset)} \mid \bigwedge \lambda(r) \leq \bigvee_{\pi(d)=r} \bigwedge (\lambda(r) \cup \rho(d)) \ (r \in R) \rangle \\ &\cong \text{Fr} \langle G \text{ (qua poset)} \mid \\ &\quad \bigwedge \lambda(r) \leq \bigvee_{D_0 \in \mathcal{F}_r D}^{\uparrow} \bigvee_{d \in D_0} \bigwedge (\lambda(r) \cup \rho(d)) \ (r \in R) \rangle \end{aligned}$$

where

$$\mathcal{F}_r D = \{D_0 \in \mathcal{F}D \mid \forall d \in D_0. \pi(d) = r\}$$

To handle the distributivity, let us write

$$\mathcal{U}(r, D_0) = \{\lambda(r) \cup \rho(d) \mid d \in D_0\}$$

Then

$$\begin{aligned} \Omega X &\cong \text{Fr}\langle G \text{ (qua poset)} \mid \\ &\quad \bigwedge \lambda(r) \leq \bigvee_{D_0 \in \mathcal{F}_r D}^\uparrow \bigwedge \{\bigvee \text{Im } \gamma \mid \gamma \in \text{Ch } \mathcal{U}(r, D_0)\} \text{ (} r \in R)\rangle \\ &\cong \text{Fr}\langle G \text{ (qua poset)} \mid \\ &\quad \bigwedge \{\bigvee V \vee g \mid g \in \lambda(r)\} \\ &\quad \leq \bigvee_{D_0 \in \mathcal{F}_r D}^\uparrow \bigwedge \{\bigvee V \vee \bigvee \text{Im } \gamma \mid \gamma \in \text{Ch } \mathcal{U}(r, D_0)\} \\ &\quad \text{(} r \in R, V \in \mathcal{F}G)\rangle \\ &\cong \text{Fr}\langle G' \text{ (qua } \vee\text{-semilattice)} \mid \\ &\quad \bigwedge \{(S \vee \{g\}) \mid g \in \lambda(r)\} \\ &\quad \leq \bigvee_{D_0 \in \mathcal{F}_r D}^\uparrow \bigwedge \{S \vee [\text{Im } \gamma] \mid \gamma \in \text{Ch } \mathcal{U}(r, D_0)\} \\ &\quad \text{(} r \in R, S \in G')\rangle \end{aligned}$$

This is what is presented by the preframe site

$$(G' \text{ (qua } \vee\text{-semilattice)}, R', D')$$

and we are done.

2.  $V \in \mathcal{F}G$  represents  $\square \bigvee_{g \in V} g \in \Omega \text{P}_U X$ . Under  $\Omega \uparrow$ , this maps to  $\bigvee V$ . Under  $\Omega \sqcap$ , it maps to  $\square \square \bigvee_{g \in V} g$ , which corresponds to  $\{\{[V]\}\}$ .

3. For  $V \in \mathcal{F}G_2$ , let us write

$$\Psi(V) = \{\phi_0 \in \mathcal{F}\phi \mid \mathcal{F} \text{snd}(\phi_0) \subseteq V\}$$

Under  $\Omega \text{P}_U f$ ,  $\square \bigvee_{h \in V} h$  maps to

$$\begin{aligned} \square \bigvee_{h \in V} \bigvee_{(S,h) \in \phi} \bigwedge S &= \square \bigvee_{\phi_0 \in \Psi(V)}^\uparrow \bigvee_{(S,g) \in \phi_0} \bigwedge S \\ &= \bigvee_{\phi_0 \in \Psi(V)}^\uparrow \square \bigwedge_{\gamma \in \text{Ch}(\{S \mid (S,g) \in \phi_0\})} \bigvee \text{Im } \gamma \\ &= \bigvee_{\phi_0 \in \Psi(V)}^\uparrow \bigwedge_{\gamma \in \text{Ch}(\{S \mid (S,g) \in \phi_0\})} \square \bigvee \text{Im } \gamma \end{aligned}$$

■

### 9. Geometricity of the double powerlocale

By combining Theorems 7.5 and 8.8 we already have :

9.1. THEOREM. *The double powerlocale is geometric.*

However, we shall gain more useful information by considering presentations for which the generators form a distributive lattice and all the relations are of the form

$$g \leq \bigvee_i^\uparrow h_i$$

with every  $h_i \leq g$ .

9.2. DEFINITION. *An ordered GRD-system  $(G \text{ (qua DL)}, R, D)$  ( $G$  being a distributive lattice, i.e. DL) is a DL-site iff it is equipped with structure and properties to make  $(G \text{ (qua } \wedge \text{-semilattice)}, R, D)$  a site and  $(G \text{ (qua } \vee \text{-semilattice)}, R, D)$  a preframe site.*

9.3. THEOREM. *Suppose a locale  $X$  is presented by a DL-site  $(G \text{ (qua DL)}, R, D)$ . Then  $\mathbb{P}X$  is presented by  $(G \text{ (qua poset)}, R, D)$ .*

PROOF. Consider the relations expressing preservation of finite joins in  $G$ :

$$\begin{aligned} \perp_G &\leq \bigvee \emptyset \\ (a \vee b) &\leq (a) \vee (b) \end{aligned}$$

By distributivity these are meet-stable. It follows that if we add them to  $R$  and  $D$  to give new sets  $R'$  and  $D'$ , then  $(G \text{ (qua } \wedge \text{-semilattice)}, R', D')$  is a site, equivalent to  $(G \text{ (qua DL)}, R, D)$ . Hence by Theorem 7.3,  $\mathbb{P}_L X$  is presented by  $(G \text{ (qua poset)}, R', D')$ , which is equivalent to  $(G \text{ (qua } \vee \text{-semilattice)}, R, D)$ . This is a preframe site, and so by Theorem 8.3 we have that  $\mathbb{P}X \cong \mathbb{P}_U \mathbb{P}_L X$  is presented by  $(G \text{ (qua poset)}, R, D)$ . ■

We now proceed to show that every presentation can be transformed into a DL-site.

9.4. LEMMA. *Let  $(G \text{ (qua } \wedge \text{-semilattice)}, R, D)$  be a site. Then it is equivalent to a DL-site  $(G' \text{ (qua DL)}, R', D')$  where*

$$\begin{aligned} G' &= \mathcal{F}G / \leq_L \\ R' &= G' \times \mathcal{F}R / \leq_L \\ \lambda_0(S, U) &= S \vee [\{\lambda_0(r) \mid r \in U\}] \\ D' &= \{(S, D_0, U) \in G' \times \mathcal{F}D \times \mathcal{F}R \mid \mathcal{F}\pi(D_0) \subseteq U\} / (\leq \times \leq_L \times \leq_L) \\ \pi(S, D_0, U) &= (S, U) \\ \rho_0(S, D_0, U) &= S \vee [\{\rho_0(d) \mid d \in D_0\}] \\ T * (S, U) &= (T \vee S, U) \\ T * (S, D_0, U) &= (T \vee S, D_0, U) \\ T \cdot (S, U) &= (T \wedge S, \{t \cdot r \mid t \in T, r \in U\}) \\ T \cdot (S, D_0, U) &= (T \wedge S, \{t \cdot d \mid t \in T, d \in D_0\}, \{t \cdot r \mid t \in T, r \in U\}) \end{aligned}$$

(Note: in expressions such as  $\{t \cdot r \mid t \in T, r \in U\}$  we have treated  $T$  and  $U$  as though they were actual sets rather than as sets modulo  $\leq_L$ . The well-definedness of the expressions is easily checked; for instance, if  $T \leq_L T'$  then  $\{t \cdot r \mid t \in T, r \in U\} \leq_L \{t \cdot r \mid t \in T', r \in U\}$ .)

PROOF. We must show that

$$\text{Fr}\langle G \text{ (qua } \wedge \text{-semilattice)} \mid \lambda_0(r) \leq \bigvee_{\pi(d)=r} \rho_0(d) \text{ (} r \in R, S \in G') \rangle$$

is isomorphic to

$$\begin{aligned} &\text{Fr}\langle \mathcal{F}G / \leq L \text{ (qua DL)} \mid \\ &\bigvee_{r \in U} S \vee \bigvee_{r \in U} \lambda_0(r) \leq \bigvee^\uparrow \{ \bigvee_{d \in D_0} S \vee \bigvee_{d \in D_0} \rho_0(d) \mid \mathcal{F}\pi(D_0) \subseteq U \} \text{ (} U \in \mathcal{F}R, S \in \mathcal{F}G) \rangle \end{aligned}$$

This is routinely proved by constructing mutually inverse homomorphisms, much as in Theorem 7.5.

It is also necessary to show that the structure described is indeed that of a DL-site. The trickiest part is  $\cdot$ . For instance, we have

$$\begin{aligned} \lambda_0(T \cdot (S, U)) &= (T \wedge S) \vee \bigvee \{ \lambda_0(t \cdot u) \mid t \in T, u \in U \} \\ &= (T \wedge S) \vee \bigvee \{ t \wedge \lambda_0(u) \mid t \in T, u \in U \} \\ &= (T \wedge S) \vee (T \wedge \bigvee \{ \lambda_0(u) \mid u \in U \}) \\ &= T \wedge \lambda_0(S, U) \end{aligned}$$

For condition 6 for a site, suppose  $\pi(S', D'_0, U') = T \cdot (S, U)$ , so  $S' = T \wedge S$  and  $U' = \{t \cdot r \mid t \in T, r \in U\}$ . For each  $d' \in D'_0$ , we have  $\pi(d') \in U'$  and so we can find  $t \in T$  and  $r \in U$  such that  $\pi(d') = t \cdot r$ , and then  $d$  such that  $\pi(d) = r$  and  $d' \leq t \cdot d$ . Hence we can find  $D_0 \in \mathcal{F}D$  such that  $\mathcal{F}\pi(D_0) \subseteq U$  and  $D'_0 \leq_L \{t \cdot d \mid t \in T, d \in D_0\}$ . Then  $\pi(S, D_0, U) = (S, U)$  and  $(S', D'_0, U') \leq T \cdot (S, D_0, U)$ . ■

9.5. THEOREM. *Every presentation can be geometrically transformed into an equivalent one given by a DL-site.*

PROOF. By Theorem 7.5, any presentation is equivalent to one given by a site; and then Lemma 9.4 it is equivalent to one given by a DL-site. ■

## 10. Geometricity of exponentiation of locales

We use the ordinary categorical definition of exponentiation.

10.1. DEFINITION. *Let  $X$  and  $Y$  be locales. Then the exponential  $Y^X$  is a locale equipped with a map  $\text{ev} : Y^X \times X \rightarrow Y$ , the evaluation map, such that the two functors  $\mathbf{Loc}(-, Y^X)$  and  $\mathbf{Loc}(- \times X, Y)$  from  $\mathbf{Loc}$  to  $\mathbf{Set}$  are naturally isomorphic under the transformation  $g \mapsto (g \times \text{Id}_X); \text{ev}$ .*

*If  $f : Z \times X \rightarrow Y$  then we write  $\ulcorner f \urcorner : Z \rightarrow Y^X$  for its corresponding exponential transpose.*

This is all completely standard. Note, however, that the natural isomorphism is necessarily an *order* isomorphism, with respect to the order enrichment  $\sqsubseteq$  (the specialization order) on  $\mathbf{Loc}$ . This is because for any locales  $Z$  and  $W$ , a pair of maps  $f, g : Z \rightarrow W$  with  $f \sqsubseteq g$  is equivalent to a map from  $\mathbb{S} \times Z \rightarrow W$ .

10.2. PROPOSITION. *Let  $X, Y$  and  $Z$  be locales. Then  $Z$  is an exponential  $Y^X$  iff there is a geometric bijection between points of  $Z$  and maps from  $X$  to  $Y$ .*

*By “geometric bijection”, we mean a bijection at every stage that is preserved by pullback along maps.*

PROOF.  $\Rightarrow$ : Let  $\text{ev} : Z \times X \rightarrow Y$  be the evaluation map. Consider stage  $W$ . Points of  $Z$  over  $W$  are equivalent to global maps  $z : W \rightarrow Z$ , and maps  $X \rightarrow Y$  over  $W$  are equivalent to global maps  $W \times X \rightarrow Y$ . By the definition of exponential there is a bijection between these, under which  $z$  corresponds to  $y = (z \times \text{Id}_X); \text{ev}$ . To show geometricity, suppose  $f : W' \rightarrow W$  is a map. Under pullback along  $f$ , we find  $z$  and  $y$  become  $f; z$  and  $(f \times \text{Id}_X); y$ , which still correspond under the bijection at  $W'$ .

$\Leftarrow$ : At stage  $Z$ , let the generic point  $\text{Id}_Z$  correspond to a map  $\text{ev} : Z \times X \rightarrow Y$ . At any stage  $W$ , geometricity tells us that maps  $z : W \rightarrow Z$  correspond to  $(z \times \text{Id}_X); \text{ev}$  as required for the exponential. ■

10.3. PROPOSITION. *Exponentiation – when it exists – is preserved by pullback.*

PROOF. This essentially follows from the previous proposition. Suppose  $Z$  is an exponential  $Y^X$  at stage  $W$  ( $X, Y$  and  $Z$  are locales over  $W$ ), and suppose  $f : W' \rightarrow W$ . For stages previous to  $W'$ , i.e. maps  $W'' \rightarrow W'$ , we have the geometric bijection between points of  $Z$  and maps  $X \rightarrow Y$ , but this amounts to a geometric bijection between points of  $f^*Z$  and maps  $f^*X \rightarrow f^*Y$ . Hence  $f^*Z$  is an exponential  $(f^*Y)^{(f^*X)}$  over  $W'$ . ■

The standard account ([Hyland 81]; or see [Johnstone 82] for a slightly different development) says –

- If  $\mathbb{S}^X$  exists then  $\Omega X$  is a continuous lattice.
- If  $\Omega X$  is a continuous lattice then  $Y^X$  exists for every locale  $Y$ .

We shall not assume these results, but our geometric development of  $\mathbb{P}$  will prove them along the way.

## 11. The first main result

In this section we shall prove the following.

11.1. THEOREM. *Let  $X$  be a locale for which the exponential  $\mathbb{S}^X$  exists. Then  $\mathbb{S}^{\mathbb{S}^X}$  also exists, and is homeomorphic to  $\mathbb{P}X$ .*

Hence  $\mathbb{S}^{(-)}$  is a kind of partial square root for  $\mathbb{P} = P_U \circ P_L$ .

Throughout the section (and the next section too), let  $X$  be a locale for which the exponential  $\mathbb{S}^X$  exists. We suppose further that  $(G_X \text{ (qua DL)}, R_X, D_X)$  and  $(G_{\mathbb{S}^X} \text{ (qua DL)}, R_{\mathbb{S}^X}, D_{\mathbb{S}^X})$  are DL-sites presenting  $X$  and  $\mathbb{S}^X$ .

11.2. DEFINITION. *We write  $\text{In} : X \times \mathbb{S}^X \rightarrow \mathbb{S}$  for the reverse of  $\text{ev} : \mathbb{S}^X \times X \rightarrow \mathbb{S}$ . We also write  $\text{In}$  for the corresponding open of  $X \times \mathbb{S}^X$  (an element of  $\Omega(X \times \mathbb{S}^X)$ ) or for the open sublocale  $\text{In} \hookrightarrow X \times \mathbb{S}^X$ .*

Since  $G_X$  and  $G_{\mathbb{S}^X}$  are bases for  $X$  and  $\mathbb{S}^X$ , the opens  $b \times V$  ( $b \in G_X$ ,  $V \in G_{\mathbb{S}^X}$ ) provide a basis for  $X \times \mathbb{S}^X$ . Hence there is a subset  $\text{In}_0$  of  $G_X \times G_{\mathbb{S}^X}$  such that  $\text{In} = \bigvee \{b \times V \mid \langle b, V \rangle \in \text{In}_0\}$ .

11.3. DEFINITION. *Let  $W \subseteq G_X \times G_{\mathbb{S}^X}$ . We say  $W$  is a finitely bilinear ideal iff*

1.  $W$  is lower closed under  $\leq \times \leq$ .
2. If  $b \in G_X$ ,  $\mathcal{V} \in \mathcal{F}G_{\mathbb{S}^X}$ , and for every  $V \in \mathcal{V}$  we have  $\langle b, V \rangle \in W$ , then  $\langle b, \bigvee \mathcal{V} \rangle \in W$ .
3. If  $U \in \mathcal{F}G_X$ ,  $V \in G_{\mathbb{S}^X}$ , and for every  $b \in U$  we have  $\langle b, V \rangle \in W$ , then  $\langle \bigvee U, V \rangle \in W$ .

11.4. LEMMA. *Given  $W \subseteq G_X \times G_{\mathbb{S}^X}$ , we can construct, intuitionistically, a finitely bilinear ideal  $\overline{W}$  such that*

$$\bigvee \{b \times V \mid \langle b, V \rangle \in \overline{W}\} = \bigvee \{b \times V \mid \langle b, V \rangle \in W\}$$

PROOF.  $\overline{W}$  is  $\{\langle b, V \rangle \in G_X \times G_{\mathbb{S}^X} \mid b \times V \leq \bigvee \{b' \times V' \mid \langle b', V' \rangle \in W\}\}$ . ■

(For present purposes, the intuitionistic construction (valid in any category of sheaves) is enough, for the property of being a finitely bilinear ideal with the required properties is geometric. However, we conjecture that  $\overline{W}$  can be constructed geometrically from  $W$ .)

It follows that, without loss of generality, we can assume that  $\text{In}_0$  is a finitely bilinear ideal.

The map  $\text{In}$  can be understood as an open in three ways. It is an open of  $X \times \mathbb{S}^X$ , but also over  $\mathbb{S}^X$  it is an open of  $X$  (which is how we shall often view it in this section), while over  $X$  it is an open of  $\mathbb{S}^X$  (which is more the viewpoint of the next section). Intuitionistic reasoning (valid in toposes) allows us to exploit these three viewpoints, but it is the more stringent geometric reasoning that allows us to move continuously between them. For instance, working over  $\mathbb{S}^X$  we view  $\text{In}$  as an open of  $X$  and can therefore find



an open  $\boxtimes$ In of  $\mathbb{P}X$ . But that is  $\mathbb{P}X$  constructed working over  $\mathbb{S}^X$ , and it is the detailed geometricity arguments that allow us to view  $\boxtimes$ In as a global open of  $\mathbb{P}X \times \mathbb{S}^X$ . We shall look at this more closely in due course.

If  $\ulcorner a \urcorner$  is a point of  $\mathbb{S}^X$ , then its exponential transpose  $a : X \rightarrow \mathbb{S}$  has  $a(x) = \text{In}\langle x, \ulcorner a \urcorner \rangle$  for each point  $x$  of  $X$ , and so we see that  $x \vDash a$  iff  $\langle x, \ulcorner a \urcorner \rangle \vDash \text{In}$ .

**11.5. PROPOSITION.** *Let  $a$  and  $U$  be opens of  $X$  and  $\mathbb{S}^X$  respectively. Then  $a \times U \leq \text{In}$  iff  $\ulcorner a \urcorner$  is a lower bound of  $U$ .*

**PROOF.**  $\Rightarrow$ : Let  $\ulcorner b \urcorner$  be a point of  $U \hookrightarrow \mathbb{S}^X$ . If  $x$  is a point of  $a$ , then  $\langle x, \ulcorner b \urcorner \rangle \vDash a \times U \leq \text{In}$  and so  $x \vDash b$ . We deduce that  $a \leq b$ , and so  $\ulcorner a \urcorner \sqsubseteq \ulcorner b \urcorner$ .

$\Leftarrow$ : Let  $\langle x, \ulcorner b \urcorner \rangle$  be a point of  $a \times U$ . Since  $\ulcorner b \urcorner \vDash U$  we have  $\ulcorner a \urcorner \sqsubseteq \ulcorner b \urcorner$ , and so  $a \leq b$ . Since  $x \vDash a$  we deduce  $x \vDash b$ , and so  $\langle x, \ulcorner b \urcorner \rangle \vDash \text{In}$ . Hence  $a \times U \leq \text{In}$ .  $\blacksquare$

It is perhaps worth pausing here to examine more closely how this very spatial argument can have any hope of yielding a valid result for locales.

First, note that the proposition is intended to be true at any stage  $Y$  – we shall certainly need this in the subsequent development. Viewed globally, the proposition then concerns opens  $a$  and  $U$  of  $Y \times X$  and  $Y \times \mathbb{S}^X$ , and the left hand condition says that, as opens of  $Y \times X \times \mathbb{S}^X$ , we have  $a \times_Y U \leq Y \times \text{In}$ . The right hand condition is a little more subtle, and we shall discuss it shortly.

Next, it suffices to prove the proposition globally, with  $Y = 1$  – just as it appears to be stated –, as long as we reason intuitionistically. This uses the geometricity of  $\mathbb{S}^X$  and its gadgetry (Proposition 10.3), which implies that

$$(Y \times \mathbb{S}^X) \times_Y (Y \times X) \cong Y \times \mathbb{S}^X \times X \xrightarrow{Y \times \text{ev}} Y \times \mathbb{S}$$

makes  $Y \times \mathbb{S}^X$  the exponential  $(Y \times \mathbb{S})^{(Y \times X)}$  over  $Y$ . It follows that the ambient assumptions about existence of  $\mathbb{S}^X$  also hold over  $Y$ , so what we prove globally about  $X$  and  $\mathbb{S}^X$  also holds (as long as we reason intuitionistically) about  $Y \times X$  and  $Y \times \mathbb{S}^X$  over  $Y$ .

Next we consider what it means to say that  $\ulcorner a \urcorner$  is a lower bound of  $U$ . Since  $U$  is an open sublocale of  $\mathbb{S}^X$ , any point of  $U$  is also a point of  $\mathbb{S}^X$  and we ask for all these to be greater than  $\ulcorner a \urcorner$ . The catch, however, is that a point of  $U$  may be at some non-global stage, and then corresponds to a point of  $\mathbb{S}^X$  at that stage. This must be compared not directly with  $\ulcorner a \urcorner$ , but with a pullback of it.

Working over  $U$ , we have a point of  $\mathbb{S}^X$  (given by the inclusion  $U \hookrightarrow \mathbb{S}^X$ ) and hence an open  $b$  of  $X$ . It is the generic open of  $X$  in  $U$ , and for any point  $u$  of  $U$  the corresponding open of  $X$  is  $u^*(b)$ . We can therefore say that  $\ulcorner a \urcorner$  is a lower bound of  $U$  iff  $a \leq b$  over  $U$ . Reviewing this globally we see that  $a$  and  $b$  correspond to two opens of  $X \times U$  and hence of  $X \times \mathbb{S}^X$ .  $a$  corresponds to  $a \times U$ , and  $b$  to  $\text{In}$ : hence we obtain the result.

We now move on to showing that each point of  $\mathbb{S}^X$  is a directed join of global points below it (in fact, way below it). However, first we prove a lemma.

11.6. LEMMA. *Let  $b$  and  $c$  be opens of  $X$  and  $1$  (so  $c$  is a point of  $\mathbb{S}$ , or a proposition).  $b \times c$  is an open of  $X \times 1 \cong X$  and so gives a point  $\ulcorner b \times c \urcorner$  of  $\mathbb{S}^X$ . Then*

$$\ulcorner b \times c \urcorner = \bigsqcup^\uparrow (\{\ulcorner \emptyset \urcorner\} \cup \{\ulcorner b \urcorner | c\})$$

PROOF. Consider exponential transposes. If  $x$  is a point of  $X$ , then

$$\langle x, \bigsqcup^\uparrow (\{\ulcorner \emptyset \urcorner\} \cup \{\ulcorner b \urcorner | c\}) \rangle = \bigsqcup^\uparrow (\{\langle x, \ulcorner \emptyset \urcorner \rangle\} \cup \{\langle x, \ulcorner b \urcorner \rangle | c\})$$

and so

$$\begin{aligned} \langle x, \bigsqcup^\uparrow (\{\ulcorner \emptyset \urcorner\} \cup \{\ulcorner b \urcorner | c\}) \rangle \vDash \text{In} &\Leftrightarrow x \vDash \emptyset \text{ or } (c \text{ and } x \vDash b) \\ &\Leftrightarrow \langle x, \ulcorner b \times c \urcorner \rangle \vDash \text{In} \end{aligned}$$

■

Our main tool is the following. It says in effect that the *generic* open of  $X$  (the generic point of  $\mathbb{S}^X$ ) is a directed join of *global* points that are in fact way below it. Clearly this is a generalization of the known fact that  $\Omega X$  is a continuous lattice, so that every *global* open of  $X$  is a directed join of globals way below it.

11.7. THEOREM. *If  $a$  is an open of  $X$  then*

$$a = \bigvee^\uparrow \{b \mid \exists V. (\langle b, V \rangle \in \text{In}_0 \text{ and } \ulcorner a \urcorner \vDash V)\}$$

PROOF. Note that the join is directed. Suppose  $\langle b_i, V_i \rangle \in \text{In}_0$  and  $\ulcorner a \urcorner \vDash V_i$  holds for  $i$  in some finite indexing set. Let  $V = \bigwedge_i V_i$ . Then  $\ulcorner a \urcorner \vDash V$  holds. Also  $\langle b_i, V \rangle \in \text{In}_0$ , so if  $b = \bigvee_i b_i$  then  $\langle b, V \rangle \in \text{In}_0$  and  $b$  is an upper bound for the  $b_i$ 's in the set.

Now for any point  $x$  of  $X$ , we have

$$\begin{aligned} x \vDash a &\Leftrightarrow \langle x, \ulcorner a \urcorner \rangle \vDash \text{In} = \bigvee \{b \times V \mid \langle b, V \rangle \in \text{In}_0\} \\ &\Leftrightarrow \exists \langle b, V \rangle \in \text{In}_0. (x \vDash b \text{ and } \ulcorner a \urcorner \vDash V) \\ &\Leftrightarrow x \vDash \bigvee^\uparrow \{b \mid \exists V. (\langle b, V \rangle \in \text{In}_0 \text{ and } \ulcorner a \urcorner \vDash V)\} \end{aligned}$$

Again, geometricity plays an important role. The result is to hold at any stage  $Y$ , but geometricity of exponentiation (Proposition 10.3) allows us to assume without loss of generality that  $Y = 1$ . After that we use Proposition 2.2. ■

11.8. COROLLARY. *If  $\ulcorner a \urcorner$  is a point of  $\mathbb{S}^X$  then*

$$\ulcorner a \urcorner = \bigsqcup^\uparrow \{\ulcorner b \urcorner \mid \exists V. (\langle b, V \rangle \in \text{In}_0 \text{ and } \ulcorner a \urcorner \vDash V)\}$$

We now turn to the role of  $\mathbb{P}X$ .

11.9. LEMMA.

1. If  $e$  is a point of  $\mathbb{P}X$  then there is an open  $U(e)$  of  $\mathbb{S}^X$  defined by

$$\ulcorner a \urcorner \vDash U(e) \text{ iff } e \vDash \boxtimes a.$$

2. If  $U$  is an open of  $\mathbb{S}^X$  then there is a point  $\ulcorner U \urcorner$  of  $\mathbb{P}X$  defined by

$$\ulcorner U \urcorner \vDash \boxtimes b \text{ iff } \ulcorner b \urcorner \vDash U$$

for  $b \in G_X$ . In fact this holds for every open  $b$  of  $X$ .

3. (1) and (2) define a geometric bijection between points of  $\mathbb{P}X$  and opens of  $\mathbb{S}^X$ .

PROOF. (1): Obvious – but relies heavily on the geometricity of the constructions of  $a$  from  $\ulcorner a \urcorner$  and of  $\boxtimes a$  from  $a$ .

(2): We must show that this definition of  $\ulcorner U \urcorner$  respects the relations used in presenting  $\mathbb{P}X$  (Theorem 9.3). Each relation is of the form  $b \leq \bigvee_i^\uparrow c_i$ , where the join is globally directed. We know, then, that  $\ulcorner b \urcorner \sqsubseteq \bigsqcup_i^\uparrow (\ulcorner c_i \urcorner)$ , so it follows that if  $\ulcorner b \urcorner \vDash U$  then  $\ulcorner c_i \urcorner \vDash U$  for some  $i$ .

The general result holds because, by Theorem 11.7, every open of  $X$  is a directed join of generating global opens from  $G_X$ .

(3): Suppose  $U$  is an open of  $\mathbb{S}^X$ . If  $\ulcorner a \urcorner$  is any point of  $\mathbb{S}^X$  then we have

$$\ulcorner a \urcorner \vDash U(\ulcorner U \urcorner) \Leftrightarrow \ulcorner U \urcorner \vDash \boxtimes a \Leftrightarrow \ulcorner a \urcorner \vDash U$$

and so  $U(\ulcorner U \urcorner) = U$ .

Now suppose  $e$  is a point of  $\mathbb{P}X$ . For any  $b \in G_X$  we then have

$$\ulcorner U(e) \urcorner \vDash \boxtimes b \Leftrightarrow \ulcorner b \urcorner \vDash U(e) \Leftrightarrow e \vDash \boxtimes b$$

and so  $\ulcorner U(e) \urcorner = e$ . ■

Applying Proposition 10.2, we have now proved Theorem 11.1. Moreover, the construction in Proposition 10.2 gives us –

11.10. PROPOSITION. *The evaluation map  $ev : \mathbb{P}X \times \mathbb{S}^X \rightarrow \mathbb{S}$  is got by working over  $\mathbb{S}^X$ . There  $In : X \times \mathbb{S}^X \rightarrow \mathbb{S}$  corresponds to an open of  $X$ . This gives an open  $\boxtimes In$  of  $\mathbb{P}X$ , and  $ev$  is the corresponding global open of  $\mathbb{P}X \times \mathbb{S}^X$ .*

Note how geometricity resolves an apparent ambiguity here: is this  $\mathbb{P}$  calculated over  $\mathbb{S}^X$ , or is it the global  $\mathbb{P}X$  pulled back to a locale  $\mathbb{P}X \times \mathbb{S}^X$  over  $\mathbb{S}^X$ ? However, the geometricity of  $\mathbb{P}$  (Theorem 9.1) gives us a canonical homeomorphism.

11.11. PROPOSITION. *As an open of  $\mathbb{P}X \times \mathbb{S}^X$ ,*

$$ev = \bigvee \{ \boxtimes b \times V \mid \langle b, V \rangle \in In_0 \}$$

PROOF. We use Theorem 11.8. Working over  $\mathbb{S}^X$  we have that  $\text{In}$  is an open of  $X$ , and the corresponding point  $\ulcorner \text{In} \urcorner$  of  $\mathbb{S}^X$  is the generic point. Hence

$$\text{In} = \bigvee^\uparrow \{b \mid \exists V. (\langle b, V \rangle \in \text{In}_0 \text{ and } V)\}$$

and

$$\text{ev} = \boxtimes \text{In} = \bigvee^\uparrow \{\boxtimes b \mid \exists V. (\langle b, V \rangle \in \text{In}_0 \text{ and } V)\}$$

because  $\boxtimes$  preserves internal directed joins. By geometricity of  $\mathbb{P}$  and its gadgetry we see that the open  $\boxtimes b$  of  $\mathbb{P}X$  over  $\mathbb{S}^X$  corresponds to the open  $\boxtimes b \times \mathbb{S}^X$  of  $\mathbb{P}X \times \mathbb{S}^X$  globally, so externally we get

$$\text{ev} = \bigvee \{(\boxtimes b \times \mathbb{S}^X) \wedge (X \times V) \mid \langle b, V \rangle \in \text{In}_0\}$$

which gives us the result. ■

## 12. $\mathbb{S}^X$ as $\mathbb{P}$ -algebra

We continue our examination of a locale  $X$  for which  $\mathbb{S}^X$  exists. We shall see that  $\mathbb{S}^X$  inherits a  $\mathbb{P}$ -algebra structure from  $\mathbb{S}$ , calculated elementwise. Moreover, the  $\mathbb{P}$ -algebra homomorphisms from  $\mathbb{S}^X$  to  $\mathbb{S}$  are equivalent to points of  $X$ . Since the global points of  $\mathbb{S}^X$  and of  $\mathbb{S}$  are the elements of  $\Omega X$  and of  $\Omega$ , this supports the contention that a  $\mathbb{P}$ -algebra is a kind of localic frame and that the  $\mathbb{P}$ -algebra structure on  $\mathbb{S}^X$  (where the exponential exists) is the localic analogue of the frame structure on  $\Omega X$ .

In this section we consider  $\text{In}$  not so much as an open of  $X$  over  $\mathbb{S}^X$ , but as an open of  $\mathbb{S}^X$  over  $X$ .

12.1. PROPOSITION. *Let  $x$  be a point of  $X$ . Under the homeomorphism  $\mathbb{P}X \cong \mathbb{S}^{\mathbb{S}^X}$ , the point  $\uparrow x$  of  $\mathbb{P}X$  corresponds to an open  $\text{In}_x = U(\uparrow x)$  of  $\mathbb{S}^X$  defined by*

$$\ulcorner a \urcorner \vDash \text{In}_x \text{ iff } \langle x, \ulcorner a \urcorner \rangle \vDash \text{In} \text{ (i.e. } x \vDash a)$$

PROOF. By Lemma 11.9,

$$\ulcorner a \urcorner \vDash U(\uparrow x) \Leftrightarrow \uparrow x \vDash \boxtimes a \Leftrightarrow x \vDash a.$$

■

The construction  $x \mapsto \text{In}_x$  is analogous to the embedding of a vector space in its double dual.

12.2. DEFINITION. *The map  $\alpha : \mathbb{P}(\mathbb{S}^X) \rightarrow \mathbb{S}^X$  is defined by*

$$\langle x, \alpha(e) \rangle \vDash \text{In} \text{ iff } e \vDash \boxtimes \text{In}_x$$

12.3. LEMMA.  $\alpha^* \text{In}_x = \boxtimes \text{In}_x$ .

PROOF.

$$e \vDash \alpha^* \text{In}_x \Leftrightarrow \alpha(e) \vDash \text{In}_x \Leftrightarrow \langle x, \alpha(e) \rangle \vDash \text{In} \Leftrightarrow e \vDash \boxtimes \text{In}_x$$

■

12.4. THEOREM.  $\alpha : \mathbb{P}(\mathbb{S}^X) \rightarrow \mathbb{S}^X$  makes  $\mathbb{S}^X$  a  $\mathbb{P}$ -algebra.

PROOF. To show  $\downarrow; \alpha = \text{Id}_{\mathbb{S}^X}$  we show

$$\langle x, \alpha \circ \downarrow(\ulcorner a \urcorner) \rangle \vDash \text{In} \Leftrightarrow \downarrow(\ulcorner a \urcorner) \vDash \boxtimes \text{In}_x \Leftrightarrow \ulcorner a \urcorner \vDash \text{In}_x \Leftrightarrow \langle x, \ulcorner a \urcorner \rangle \vDash \text{In}$$

so  $\alpha \circ \downarrow(\ulcorner a \urcorner) = \ulcorner a \urcorner$ .

Similarly, for  $\mathbb{P}\alpha; \alpha = \text{H}; \alpha$ .

$$\begin{aligned} \langle x, \alpha \circ \mathbb{P}\alpha(\zeta) \rangle \vDash \text{In} &\Leftrightarrow \mathbb{P}\alpha(\zeta) \vDash \boxtimes \text{In}_x \\ &\Leftrightarrow \zeta \vDash \boxtimes \boxtimes \text{In}_x \text{ (using Lemma 12.3)} \\ &\Leftrightarrow \text{H}(\zeta) \vDash \boxtimes \text{In}_x \Leftrightarrow \langle x, \alpha \circ \text{H}(\zeta) \rangle \vDash \text{In} \end{aligned}$$

■

12.5. PROPOSITION. *If  $x$  is a point of  $X$ , then the map  $\text{In}_x : \mathbb{S}^X \rightarrow \mathbb{S}$  is a  $\mathbb{P}$ -algebra homomorphism.*

PROOF. By definition of  $\alpha$ , the composite  $\alpha; \text{In}_x$  corresponds to the open  $\boxtimes \text{In}_x$  of  $\mathbb{P}(\mathbb{S}^X)$ , and this is just  $\mathbb{P}\text{In}_x; \boxtimes\{\top\}$ . Since  $\boxtimes\{\top\}$  is the  $\mathbb{P}$ -algebra structure map for  $\mathbb{S}$ , this is enough to show the result. ■

12.6. THEOREM. *The transformation  $x \mapsto \text{In}_x$  gives a bijection between points  $x$  of  $X$  and  $\mathbb{P}$ -algebra homomorphisms from  $\mathbb{S}^X$  to  $\mathbb{S}$ .*

PROOF. By Theorem 11.1 we have a bijection between points  $\ulcorner U \urcorner$  of  $\mathbb{P}X$  and opens  $U$  of  $\mathbb{S}^X$ , and from Proposition 12.1 we see that if  $\ulcorner U \urcorner$  can be expressed as  $\downarrow(x)$  ( $x$  then is unique, since  $\downarrow$  is an inclusion) then  $U = \text{In}_x$ . Therefore, it suffices to show that  $\ulcorner U \urcorner$  is in the image of  $\downarrow$  iff  $U : \mathbb{S}^X \rightarrow \mathbb{S}$  is a  $\mathbb{P}$ -algebra homomorphism and the  $\Rightarrow$  direction has already been done.

Suppose, then, that  $U$  is an open of  $\mathbb{S}^X$  such that  $U : \mathbb{S}^X \rightarrow \mathbb{S}$  is a  $\mathbb{P}$ -algebra homomorphism. The point  $\ulcorner U \urcorner$  of  $\mathbb{P}X$  is defined by

$$\ulcorner U \urcorner \vDash \boxtimes b \text{ iff } \ulcorner b \urcorner \vDash U$$

for  $b \in G_X$ . We show that  $\ulcorner U \urcorner$  is in the image of  $\downarrow$ . If  $\ulcorner U \urcorner = \downarrow(x)$  then  $x$  must be defined by

$$x \vDash b \text{ iff } \ulcorner b \urcorner \vDash U$$

and so what we are required to show is that this respects the lattice structure of  $G_X$ . This is clear, because from the presentation of  $X$  we know that the map  $G_X \rightarrow \mathbb{S}^X$ ,  $b \mapsto \ulcorner b \urcorner$ , is a (localic) lattice homomorphism, and so also is  $U : \mathbb{S}^X \rightarrow \mathbb{S}$  because it is a  $\mathbb{P}$ -homomorphism. ■

12.7. THEOREM. *Let  $X$  be a locale for which  $\mathbb{S}^X$  exists, and let  $W$  be a locale such that  $\mathbb{P}W$  is homeomorphic to  $\mathbb{S}^X$ . Then  $\mathbb{S}^W$  exists and is homeomorphic to  $X$ .*

PROOF. Opens of  $W$ , maps  $W \rightarrow \mathbb{S}$ , are equivalent to  $\mathbb{P}$ -algebra homomorphisms  $\mathbb{P}W \rightarrow \mathbb{S}$ . The exponential  $\mathbb{S}^X$  is also a  $\mathbb{P}$ -algebra, and so, by Corollary 4.6, its homeomorphism with  $\mathbb{P}W$  is a  $\mathbb{P}$ -algebra isomorphism. Thus the  $\mathbb{P}$ -algebra homomorphisms  $\mathbb{P}W \rightarrow \mathbb{S}$  are in bijection with the  $\mathbb{P}$ -algebra homomorphisms  $\mathbb{S}^X \rightarrow \mathbb{S}$ , and by Theorem 12.6 these are in bijection with the points of  $X$ . Hence  $X$  serves as an exponential  $\mathbb{S}^W$ . ■

### 13. Characterizing exponentiability

We have mostly been studying the consequences of the existence of  $\mathbb{S}^X$ , and where (in Theorem 12.7) we gave a sufficient condition, it relied on the existence already established of another exponential. In order to get off the ground we must find some locales  $X$  for which  $\mathbb{S}^X$  is known to exist. It suffices to use the continuous dcpos.

We shall give a somewhat indirect proof of what is, after all, a well-known result: that continuous dcpos are exponentiable as locales. When continuous dcpos are understood particularly as ideal completions of continuous information systems, it can be deduced from the following Theorem, which is intuitionistically valid.

13.1. THEOREM. *Let  $P$  be a continuous information system. Then  $\Omega \text{Idl}(P)$  is isomorphic to the lattice of rounded upper subsets of  $P$ .*

PROOF. [Vickers 93] Note that the result is not geometric, as it is not even stated geometrically. ■

However, to avoid the explicit use of frames we shall give a geometric proof. We first prove a special case.

13.2. DEFINITION. [Smyth 77] *An R-structure is a continuous information system  $P$  in which, for every  $s \in P$ ,  $\downarrow s = \{t \in P \mid t < s\}$  is an ideal.*

13.3. LEMMA. *Let  $P$  be an R-structure, let  $X = \text{Idl } P$  and let  $Y = \text{Idl}(\mathcal{F}P, (<_U)^{\text{op}})$ . Then  $Y$  serves as the exponential  $\mathbb{S}^X$ .*

PROOF. If  $a$  is an open of  $X$ , then we can define a point  $J(a)$  of  $Y$  by

$$J(a) = \{S \in \mathcal{F}P \mid \forall s \in S. \downarrow s \vDash a\}.$$

Since

$$\downarrow s = \bigsqcup_{s' < s} \uparrow \downarrow s',$$

we find that if  $\downarrow s \vDash a$  then  $\downarrow s' \vDash a$  for some  $s' < s$ . It follows that if  $S \in J(a)$  then  $S' \in J(a)$  for some  $S' <_U S$  ( $J(a)$  is rounded). If  $S_i \in J(a)$  ( $1 \leq i \leq n$ ) then we can find  $S'_i \in J(a)$  with  $S'_i <_U S_i$ . Taking  $S = \bigcup_i S'_i$ , we get that  $S \in J(a)$  and  $S <_U S_i$ . After this it is easy to conclude that  $J(a)$  is an ideal.

Now suppose  $J$  is a point of  $Y$ . We can define an open  $a(J)$  of  $X$  by

$$I \vDash a(J) \text{ iff } \exists s \in I. \{s\} \in J$$

Starting from  $J$ , we have  $J = J(a(J))$ . For

$$\begin{aligned} S \in J(a(J)) &\Leftrightarrow \forall s \in S. \{s\} \in J \\ &\Leftrightarrow \exists S' \in J. \forall s \in S. S' <_U \{s\} \\ &\Leftrightarrow \exists S' \in J. S' <_U S \Leftrightarrow S \in J \end{aligned}$$

Starting from  $a$ , we have  $a = a(J(a))$ . For if  $I$  is a point of  $X$  then  $I = \bigsqcup_{s \in I}^\uparrow \downarrow s$  and so

$$I \vDash a(J(a)) \Leftrightarrow \exists s \in I. \downarrow s \vDash a \Leftrightarrow I \vDash a$$

Since all these constructions are geometric, we can now call on Proposition 10.2.  $\blacksquare$

**13.4. THEOREM.** *Let  $P$  be a continuous information system, let  $W = \text{Idl } P$  and let  $X = \text{Idl}(\mathcal{F}P, (<_U)^{\text{op}})$ . Then  $X$  serves as the exponential  $\mathbb{S}^W$ .*

**PROOF.**  $(\mathcal{F}P, (<_U)^{\text{op}})$  is an R-structure. For suppose  $S <_U T_i$  ( $1 \leq i \leq n$ ). For each  $i$  we can find  $T'_i$  with  $S <_U T'_i <_U T_i$ , and then  $S <_U \bigcup_i T'_i <_U T_i$ . It follows by Lemma 13.3 that  $\mathbb{S}^X$  exists and is  $\text{Idl}(\mathcal{F}\mathcal{F}P, \prec)$ , where

$$\mathcal{U} \prec \mathcal{V} \text{ iff } \forall U \in \mathcal{U}. \exists V \in \mathcal{V}. U <_U V$$

By Theorem 3.11 we see that this is  $\text{P}_L(\text{P}_U(\text{Idl } X))$ , i.e.  $\mathbb{P}(\text{Idl } X)$ , and it remains to apply Theorem 12.7.  $\blacksquare$

**13.5. THEOREM.** *Let  $X$  be a locale for which  $\mathbb{S}^X$  exists, with lattices of generators  $G_X$  and  $G_{\mathbb{S}^X}$  and so on as in Section 11. We define the relation  $<$  on  $G_X$  by*

$$b < a \text{ iff } \exists V. (\langle b, V \rangle \in \text{In}_0 \text{ and } \ulcorner a \urcorner \vDash V).$$

*Then  $<$  makes  $G_X$  a continuous information system, and  $\mathbb{S}^X$  is homeomorphic to  $\text{Idl}(G_X)$ .*

**PROOF.** First we must show that  $<$  is transitive and interpolative. It will be helpful later to generalize  $b < a$  slightly, to allow  $a$  to be an arbitrary open of  $X$ . We therefore define

$$I_a = \{b \in G_X \mid \exists V. (\langle b, V \rangle \in \text{In}_0 \text{ and } \ulcorner a \urcorner \vDash V)\}$$

so that, when  $a$  and  $b$  are both in  $G_X$ , we have  $b < a$  iff  $b \in I_a$ .  $I_a$  is lower closed with respect to  $<$ , for suppose  $c < b \in I_a$ , with  $\langle c, W \rangle, \langle b, V \rangle \in \text{In}_0$ ,  $\ulcorner b \urcorner \vDash W$  and  $\ulcorner a \urcorner \vDash V$ . By Proposition 11.5,  $\ulcorner c \urcorner$  is a lower bound of  $W$  and so  $c \leq b$ . Hence  $\langle c, V \rangle \in \text{In}_0$  and we have  $c \in I_a$ . It follows that  $<$  is transitive. Next we show that  $I_a$  is rounded with respect to  $<$ . Suppose  $c \in I_a$ , with  $\langle c, W \rangle \in \text{In}_0$  and  $\ulcorner a \urcorner \vDash W$ . By Theorem 11.8 we know that

$\lceil a \rceil = \bigsqcup^\uparrow \{\lceil b \rceil \mid b \in I_a\}$ , so it follows that we have some  $b \in I_a$  with  $\lceil b \rceil \vDash W$ , so  $c < b$ . It follows that  $<$  is interpolative.

In fact,  $I_a$  is an ideal with respect to  $<$ , as we now show. Suppose  $c_i \in I_a$  for  $i$  in some finite indexing set, with  $\langle c_i, W_i \rangle \in \text{In}_0$  and  $\lceil a \rceil \vDash W_i$ . Then  $\lceil a \rceil \vDash \bigwedge_i W_i$ ,  $\langle c_i, \bigwedge_i W_i \rangle \in \text{In}_0$  and so  $\langle \bigvee_i c_i, \bigwedge_i W_i \rangle \in \text{In}_0$ . It follows that  $\bigvee_i c_i \in I_a$ . Now let us find  $b$  such that  $\bigvee_i c_i < b \in I_a$ , with  $\langle \bigvee_i c_i, V \rangle \in \text{In}_0$  and  $\lceil b \rceil \vDash V$ . We have  $\langle c_i, V \rangle \in \text{In}_0$  and it follows that  $c_i < b$ .

We have now shown that for each open  $a$  we can define an ideal  $I_a$  such that  $a$  is the directed join of the opens in  $I_a$ . It remains only to show that if  $I$  is an ideal and  $a$  is the join of its elements, then  $I = I_a$ . First, if  $c \in I$ , then we can find  $b$  such that  $c < b \in I$ . We have  $b \leq a$ , and it follows that  $c \in I_b \subseteq I_a$ . For the converse, suppose  $c \in I_a$  with  $\langle c, W \rangle \in \text{In}_0$  and  $\lceil a \rceil \vDash W$ . By Scott openness of  $W$ , and the fact that  $a$  is a *directed* join of the elements of  $I$ , it follows that there is some  $b \in I$  such that  $\lceil b \rceil \vDash W$ . Then  $c < b$ , so  $c \in I$ .

The geometricity of this argument suffices to show that  $\mathbb{S}^X$  is homeomorphic to  $\text{Idl}(G_X)$ . ■

**13.6. COROLLARY.** *Let  $X$  be a locale for which  $\mathbb{S}^X$  exists. Then  $\Omega X$  is a continuous lattice, and  $\Omega \mathbb{S}^X$  is the Scott topology on  $\Omega X$ . Its way below relation  $\ll$  satisfies*

$$b \ll a \text{ iff } \exists V \in \Omega \mathbb{S}^X. (b \times V \leq \text{In} \text{ and } \lceil a \rceil \vDash V)$$

We can now characterize those locales  $X$  for which  $\mathbb{S}^X$  exists. Note that the characterization does not require a universal quantification over all locales, but simply the existence of certain structure and properties. In this respect it is like the well known characterization by local compactness, that  $\Omega X$  is a continuous lattice. However, we use structure and properties that are geometric.

**13.7. THEOREM.** *Let  $W$  be a locale. Then  $\mathbb{S}^W$  exists iff there is a continuous information system  $X$  such that  $\mathbb{P}W$  is homeomorphic to  $\text{Idl}(\mathcal{F}X, (<_U)^{\text{op}})$ , and in that case  $\mathbb{S}^W$  is homeomorphic to  $\text{Idl}(X)$ .*

**PROOF.**  $\Leftarrow$ : Use Theorem 12.7.

$\Rightarrow$ : Let  $X$  be the information system  $(G_W, <)$  defined as in Theorem 13.5. Then  $\mathbb{S}^W$  is homeomorphic to  $\text{Idl}(X)$  and so

$$\mathbb{P}W \cong \mathbb{S}^{\mathbb{S}^W} \cong \mathbb{S}^{\text{Idl}(X)} \cong \text{Idl}(\mathcal{F}X, (<_U)^{\text{op}})$$

■

To round off this section, we give a geometric proof to show how exponentiability of  $X$  can be deduced from the existence of  $\mathbb{S}^X$ .

**13.8. LEMMA.** *Any discrete locale is exponentiable.*



PROOF. Let  $S$  be a discrete locale (i.e. a set). Let  $Y$  be an arbitrary locale, with presentation  $(G, R, D)$ , and let  $Y'$  be presented by  $(G \times S, R \times S, D \times S)$  with

$$\begin{aligned} \pi(\langle d, s \rangle) &= \langle \pi(d), s \rangle \\ \lambda(\langle r, s \rangle) &= \{ \langle g, s \rangle \mid g \in \lambda(r) \} \\ \rho(\langle d, s \rangle) &= \{ \langle g, s \rangle \mid g \in \rho(d) \} \end{aligned}$$

If  $f : S \rightarrow Y$  then we can define a point  $f'$  of  $Y'$  by

$$f' \vDash \langle g, s \rangle \text{ iff } f(s) \vDash g$$

and, conversely, for any point  $f'$  of  $Y'$  the same equation defines a map  $f : S \rightarrow Y$ .

This defines a geometric bijection between points of  $Y'$  and maps  $S \rightarrow Y$ . ■

13.9. THEOREM. *Let  $X$  be a locale for which  $\mathbb{S}^X$  exists. Then  $X$  is exponentiable.*

PROOF. Let  $Y$  be a locale, presented by  $(G, R, D)$ . Since  $\mathbb{S}^X$  exists, we see by the lemma that  $(\mathbb{S}^X)^G$  does too, and we then have

$$(\mathbb{S}^G)^X \cong \mathbb{S}^{X \times G} \cong (\mathbb{S}^X)^G$$

and, similarly,

$$(\mathbb{S}^R)^X \cong \mathbb{S}^{X \times R} \cong (\mathbb{S}^X)^R$$

Consider two maps  $\phi_\lambda, \phi_\rho : \mathbb{S}^G \rightarrow \mathbb{S}^R$  defined as follows.  $\phi_\lambda$  is the exponential transpose of the open  $\phi'_\lambda$  of  $\mathbb{S}^G \times R$ ,

$$\langle \ulcorner a \urcorner, r \rangle \vDash \phi'_\lambda \text{ iff } \forall g \in \lambda(r). g \vDash a$$

$\phi'_\rho$  is the exponential transpose of the open  $\phi'_\rho$  of  $\mathbb{S}^G \times R$ ,

$$\langle \ulcorner a \urcorner, r \rangle \vDash \phi'_\rho \text{ iff } \exists d. (\pi(d) = r \text{ and } \forall g \in \rho(d). g \vDash a)$$

From these we define two maps  $\psi_\lambda$  and  $\psi_\rho$  from  $(\mathbb{S}^G)^X$  to  $(\mathbb{S}^R)^X$ , by  $\psi_\lambda = \phi_\lambda^X$  and  $\psi_\rho = \phi_\rho^X$ .

Let  $E \hookrightarrow (\mathbb{S}^G)^X$  be the inserter from  $\psi_\lambda$  to  $\psi_\rho$ , i.e. the greatest sublocale restricted to which we have  $\psi_\lambda \sqsubseteq \psi_\rho$ . Then  $E$  is the exponential  $Y^X$ . For if  $f : X \rightarrow Y$  then we can define an open  $f'$  of  $X \times G$  (corresponding to a point of  $\mathbb{S}^{X \times G}$ ) by

$$\langle x, g \rangle \vDash f' \text{ iff } f(x) \vDash g$$

and the relations for  $Y$ , holding for each point  $f(x)$ , tell us that  $f'$  is in  $E$ . Conversely, if  $f'$  is an open of  $X \times G$  that happens to be in  $E$  then the same definition in reverse defines a point  $f(x)$  of  $Y$  for each  $x$ .

This gives a geometric bijection between points of  $E$  and maps  $X \rightarrow Y$ . ■

## 14. Conclusion

The technical results proved here seem rather specialized. We have taken the upper and lower powerlocales, known principally from computer science, and composed them to give a hitherto unstudied “double powerlocale”  $\mathbb{P}$ . This is then related to exponentiation by the homeomorphism  $\mathbb{P}X \cong \mathbb{S}^{\mathbb{S}^X}$ , provided  $\mathbb{S}^X$  exists.

$\mathbb{P}$  is the functor part of a monad on **Loc**, and its algebras bear some analogies with frames, but carried by locales instead of sets. Following through this analogy, we find that  $\mathbb{S}^X$  inherits  $\mathbb{P}$ -algebra structure from  $\mathbb{S} \cong \mathbb{P}\emptyset$ , and  $\mathbb{P}$ -algebra homomorphisms from  $\mathbb{S}^X$  to  $\mathbb{S}$  are equivalent to points of  $X$ . We deduce that if  $\mathbb{S}^X \cong \mathbb{P}W$  then  $X \cong \mathbb{S}^W$ . An intriguing direction of further study is that of  $(\mathbb{P}\text{-alg})^{\text{op}}$  as a category **Coloc** of “localic locales” or *colocales*.

In addition to investigating  $\mathbb{P}$  we have also taken the opportunity to develop techniques of *geometricity* – that is, structure and properties preserved by pullback along maps (geometric morphisms). It is these techniques that make possible a much more spatial development of locale theory using *points* of locales. For these points must be generalized points, i.e. arbitrary maps targeted on the locale. It is the geometricity that allows us to cope continuously with the variation of source. Of particular importance is the geometricity of the powerlocale constructions themselves. The treatment relies on the fact that though frames are not geometric representations of the locales (because frame structure is not preserved by inverse image functors), frame presentations *are* geometric.

Our primary application of the long section of geometricity results is thus to justify a spatial treatment of locales. However, an additional hope is that they might contribute to the foundations of constructive topology. In particular, we have in mind the *formal topologies* of [Sambin 87]. He follows Martin-Löf’s school of constructive type theory in rejecting impredicative arguments, and this includes many arguments that are valid in (the categories of sheaves over) toposes. (On the other hand, that school admits more of the axiom choice than is valid in toposes.) This leads to a rather fundamental contrast between formal topology and locale theory as conducted in toposes: although they both treat opens as more fundamental than points, formal topology cannot admit the frames as algebras carried by sets. In effect various kinds of formal topology are various ways of *presenting* frames (sometimes with extra structure, such as positivity predicates when presenting open locales). Since our geometricity requirements lead us also to reject the frames, and to develop ways of handling locales without the frames, we hope that our methods will also find application in formal topology and help to find common ground between the two approaches to constructive topology.

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